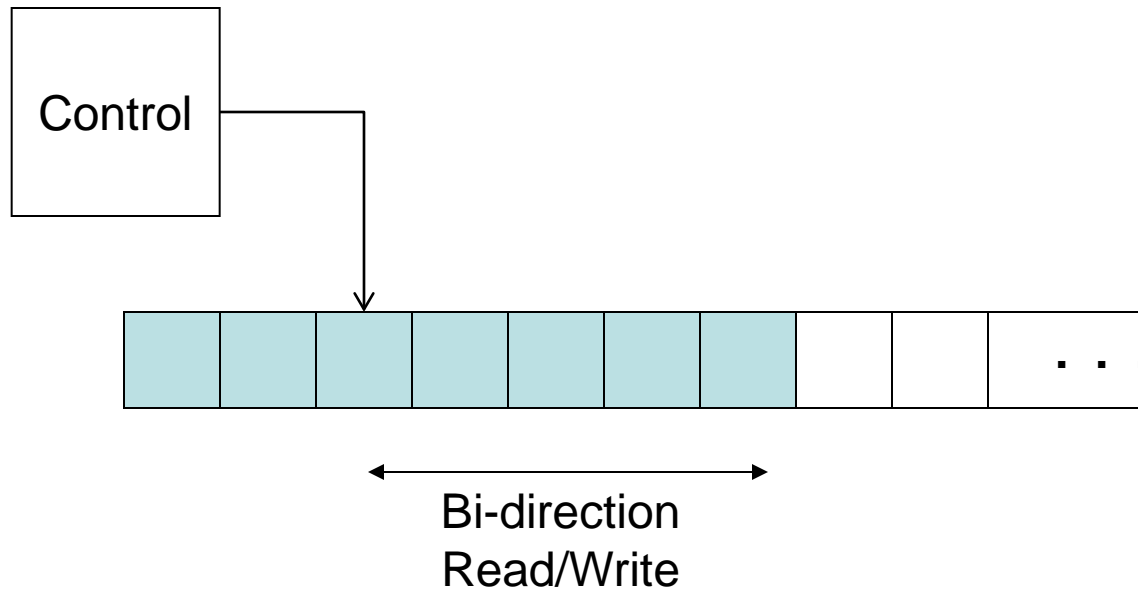


# Chapter 3: The Church-Turing Thesis

# Turing Machine (TM)



Turing machine is a much more powerful model, proposed by Alan Turing in 1936.

# Church/Turing Thesis

Anything that an algorithm can compute can be computed by a Turing machine and vice-versa.

# Turing Machine (cont'd)

- The differences between finite automata and Turing machines
  - A Turing machine can both write on the tape and read from it,
  - The read-write head can move both to the left and to the right,
  - The tape is infinite,
  - The special states for rejecting and accepting take effects immediately.

# Turing Machine Definition

A deterministic Turing machine is a 7-tuple:

$(Q, \Sigma, \Gamma, \delta, q_0, q_{\text{accept}}, q_{\text{reject}})$ , where  $Q, \Sigma, \Gamma$  are finite sets, and

1.  $Q$ : set of states

2.  $\Sigma$ : input alphabet not containing the blank symbol  $\sqcup$

3.  $\Gamma$ : tape alphabet, including  $\Sigma$  and the blank symbol

4.  $\delta: Q \times \Gamma \longrightarrow Q \times \Gamma \times \{L, R\}$  is the transition function

5.  $q_0: q_0 \in Q$  is the start state,

6.  $q_{\text{accept}}: q_{\text{accept}} \in Q$  is the accept state

7.  $q_{\text{reject}}: q_{\text{reject}} \in Q$  is the reject state

# Non-Deterministic Turing Machine

A non-deterministic Turing machine is a 7-tuple:  
( $Q, \Sigma, \Gamma, \delta, q_0, q_{\text{accept}}, q_{\text{reject}}$ ), where  $Q, \Sigma, \Gamma$  are finite sets,  
and

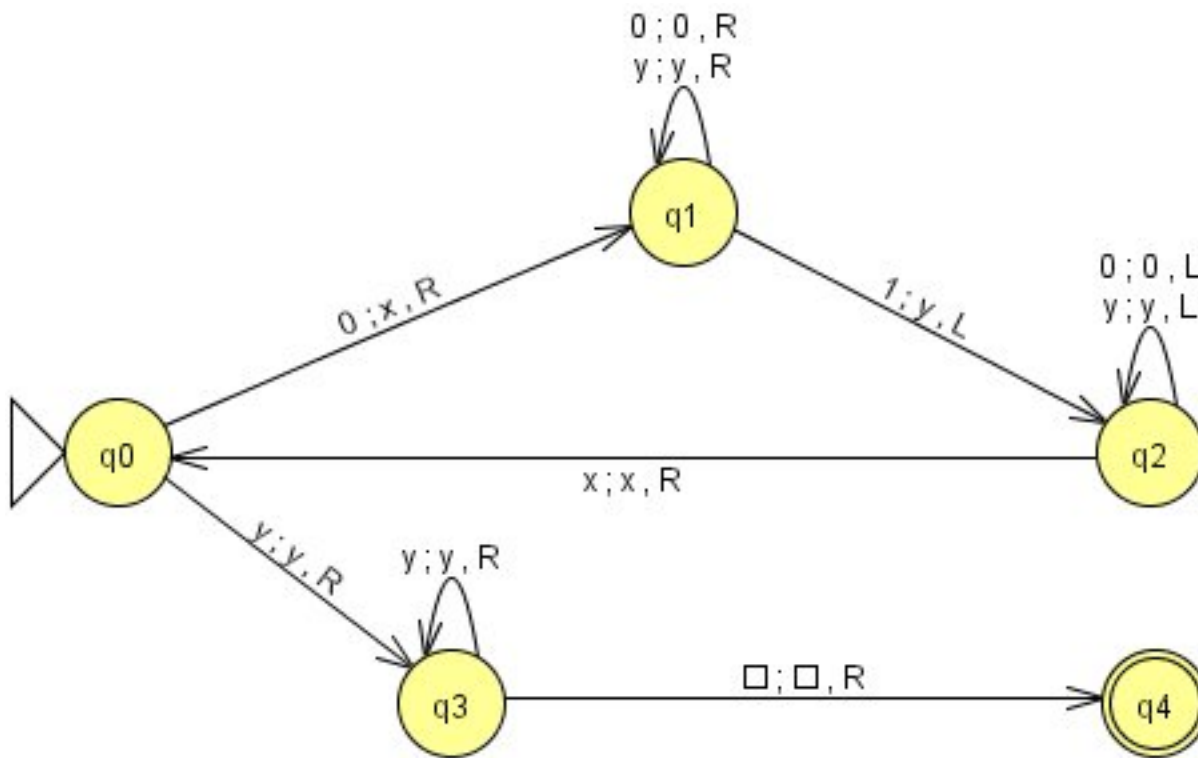
1.  $Q$ : set of states
2.  $\Sigma$ : input alphabet not containing the blank symbol  
 $\sqcup$
3.  $\Gamma$ : tape alphabet, including  $\Sigma$  and the blank symbol
4.  $\delta: Q \times \Gamma \longrightarrow P(Q \times \Gamma \times \{L, R\})$  is the transition function
5.  $q_0: q_0 \in Q$  is the start state,
6.  $q_{\text{accept}}: q_{\text{accept}} \in Q$  is the accept state
7.  $q_{\text{reject}}: q_{\text{reject}} \in Q$  is the reject state

# TM for $\{0^n 1^n \mid n > 0\}$

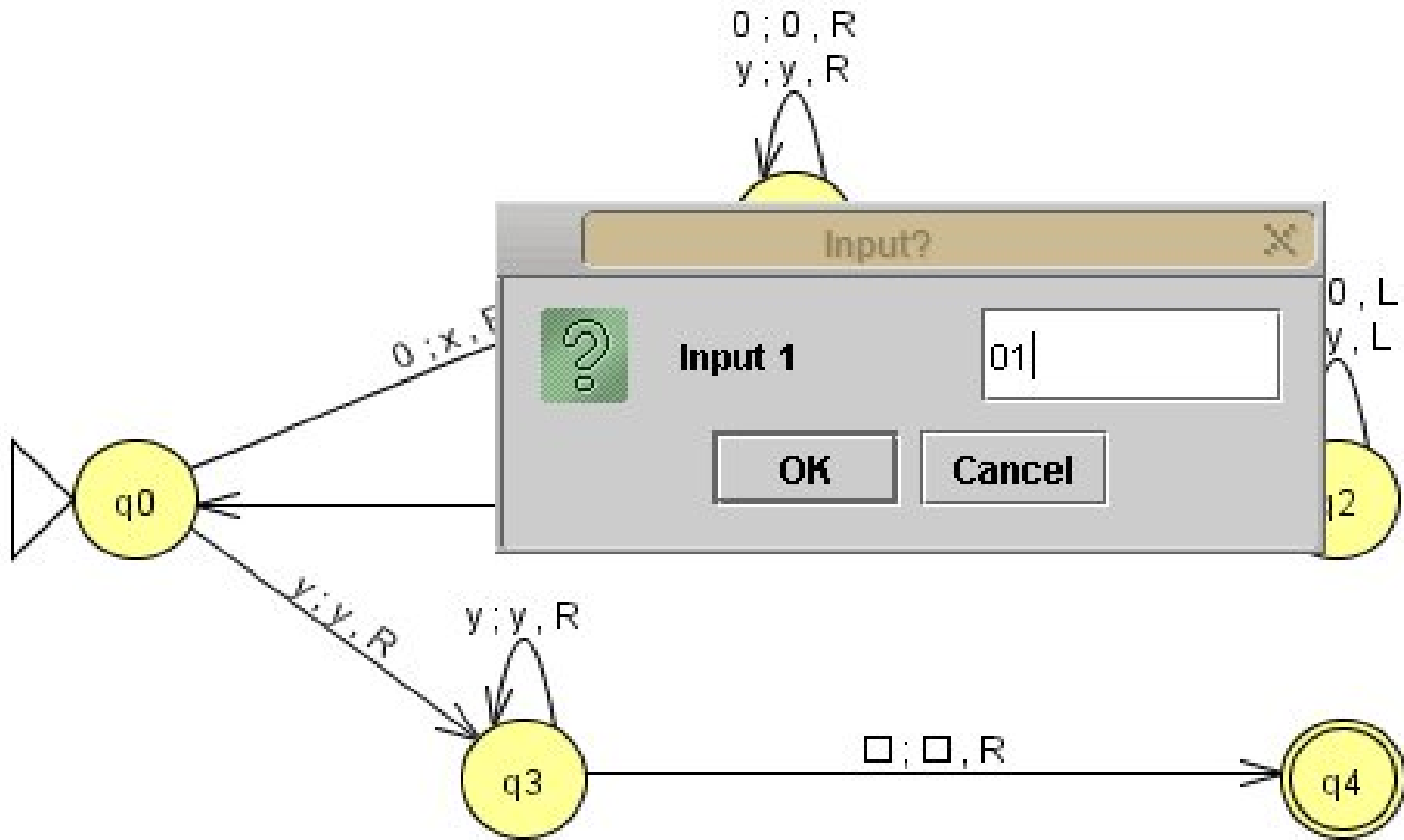
M= “on input string w:

1. Look for 0's from the left end of the tape. Only blank cells are allowed to pass.
2. If 0 found, change it to x and move right, else reject
3. Scan (to the right) passing 0's and y's until reach 1
4. If 1 found, change it to y and move left, else reject.
5. Move left passing y's and 0's
6. If x found move right
7. If 0 found, loop back to step 2.
8. If 0 not found, scan to the right passing y's and accept; otherwise reject.”

# Example of TM for $\{0^n 1^n \mid n > 0\}$



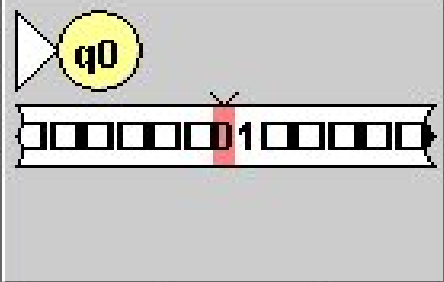
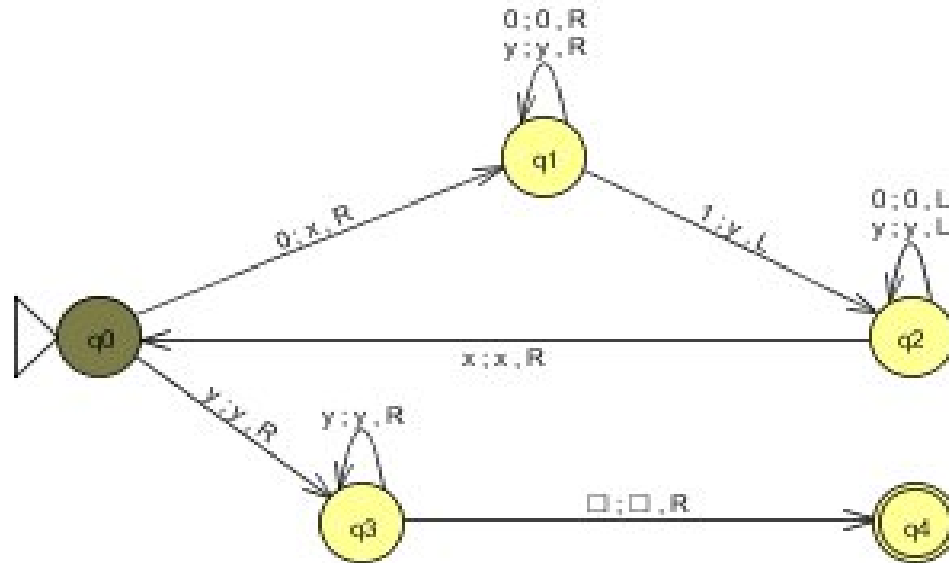




File Input Test Convert Help

Editor

Simulate: [01]

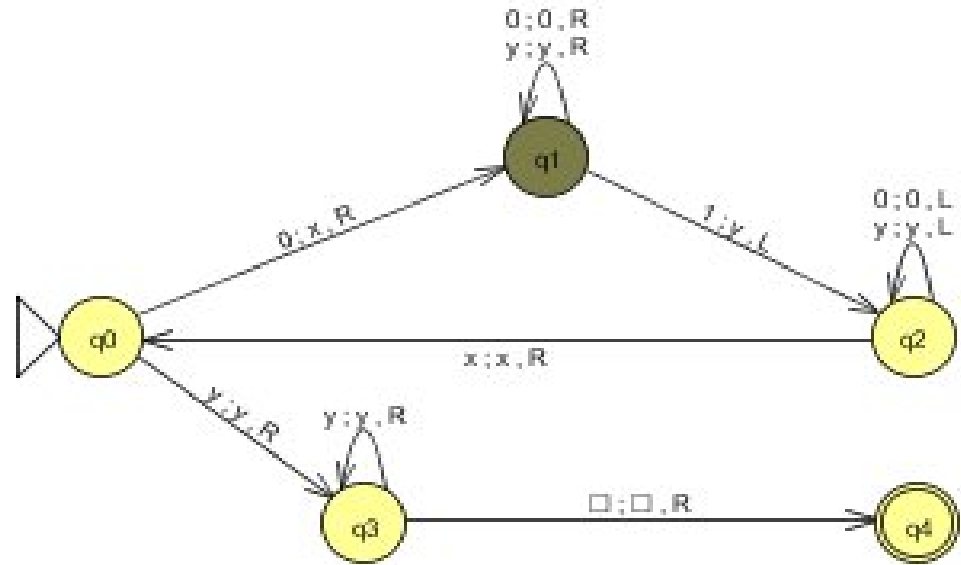


Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor

Simulate: [01]



q1

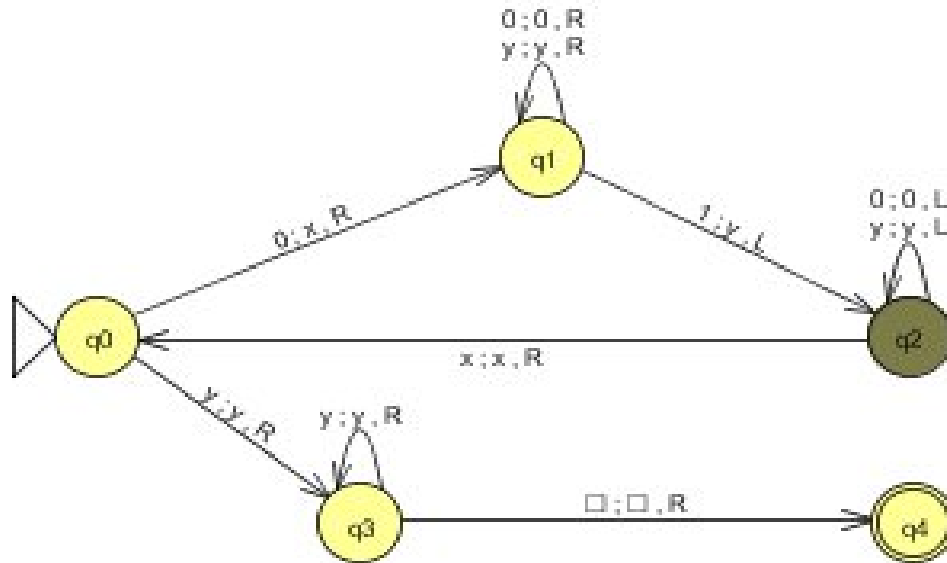


Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor

Simulate: [01]



q2



Step

Reset

Freeze

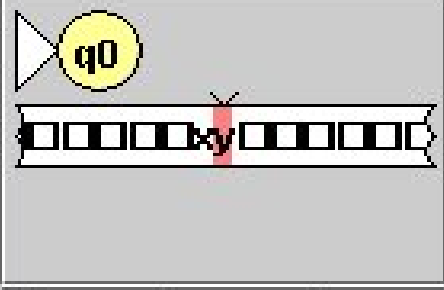
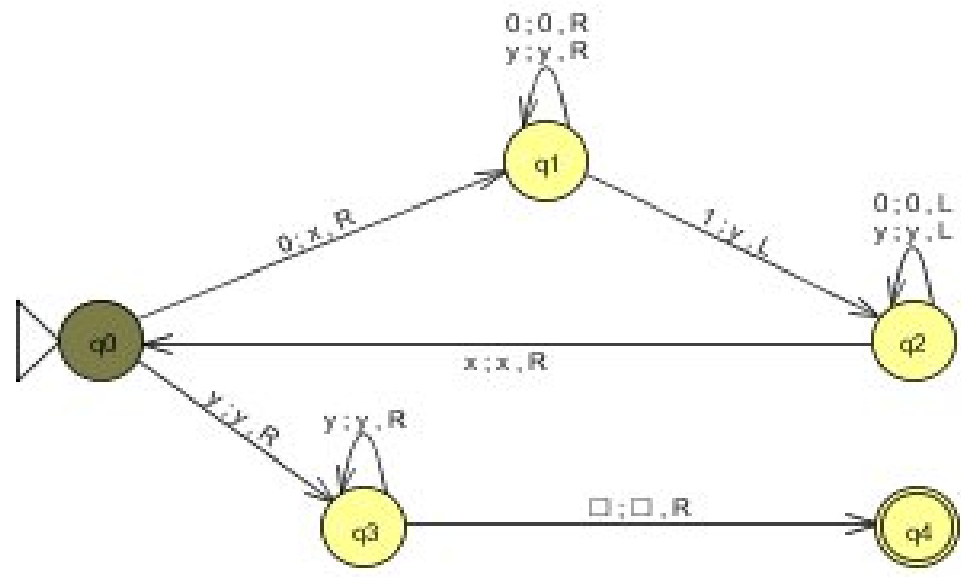
Thaw

Trace

Remove

File Input Test Convert Help

Editor Simulate: [01]

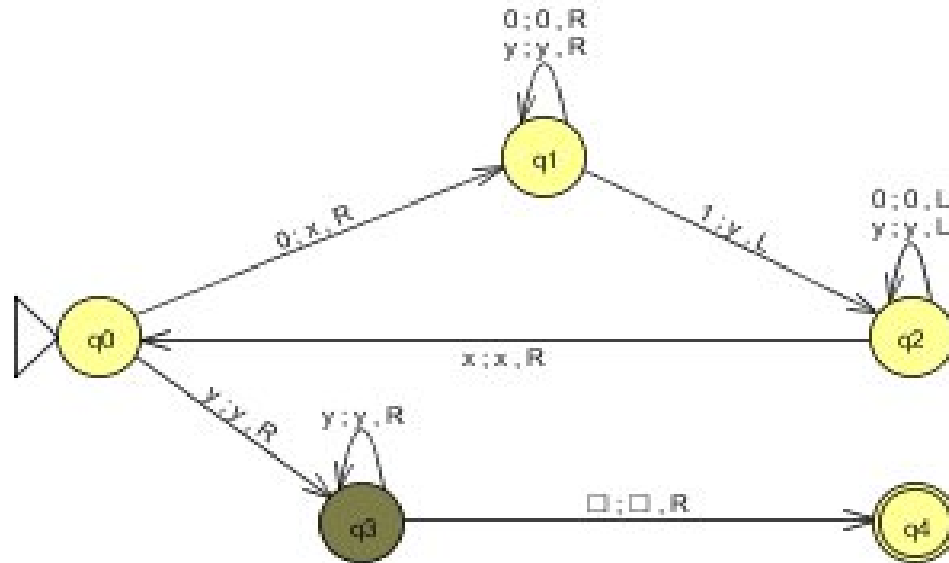


Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor

Simulate: [01]



q3



Step

Reset

Freeze

Thaw

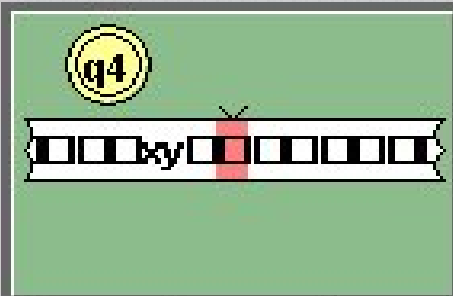
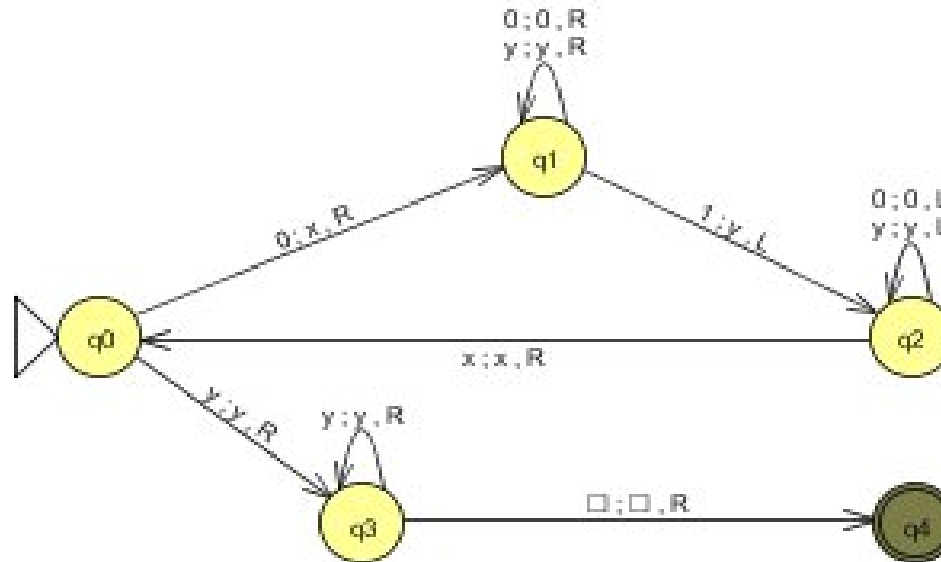
Trace

Remove

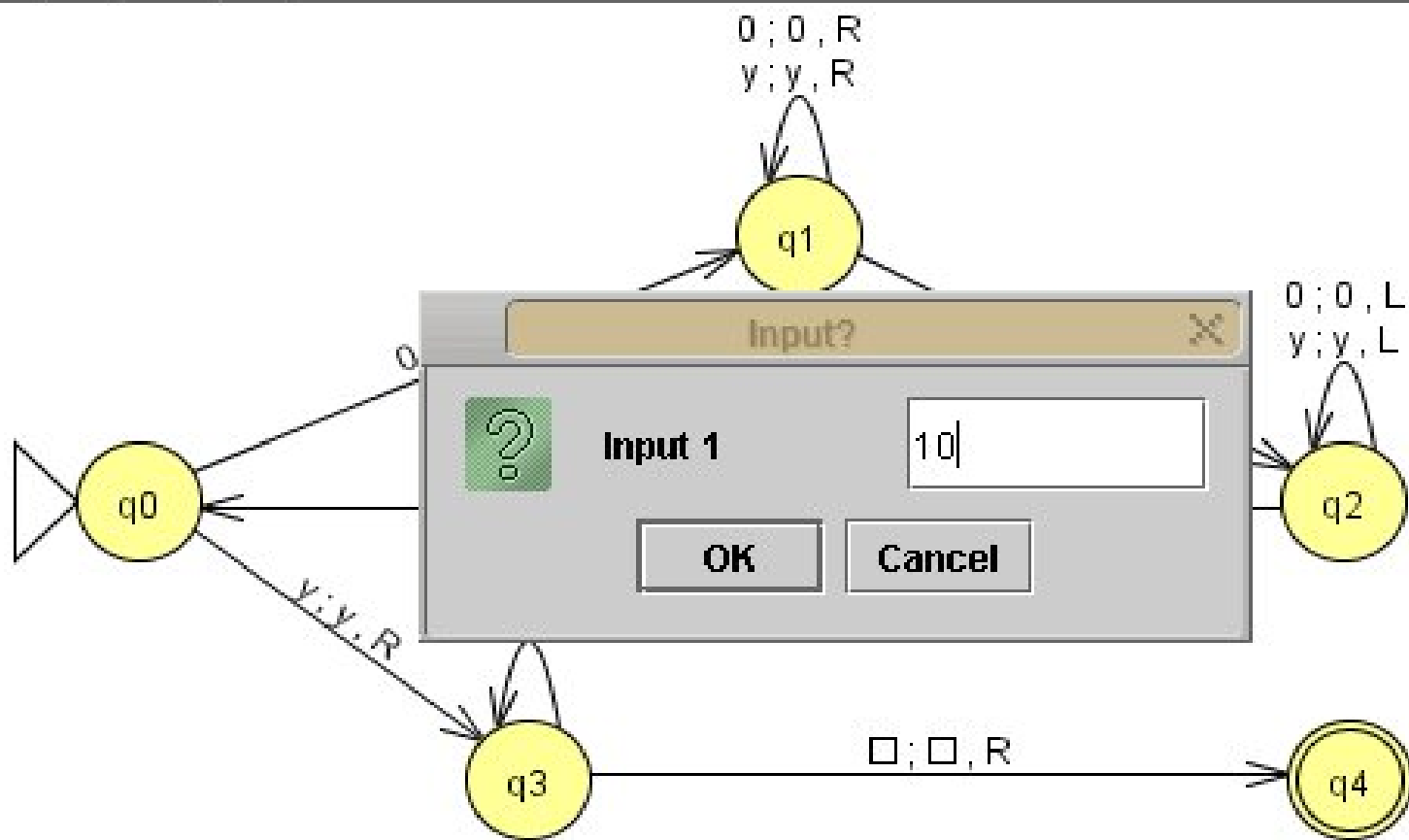
File Input Test Convert Help

Editor

Simulate: [01]



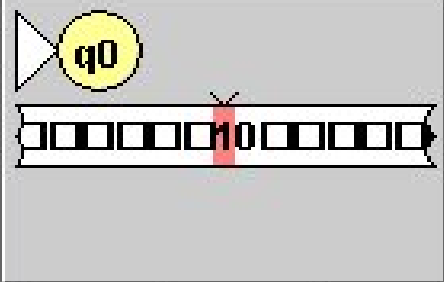
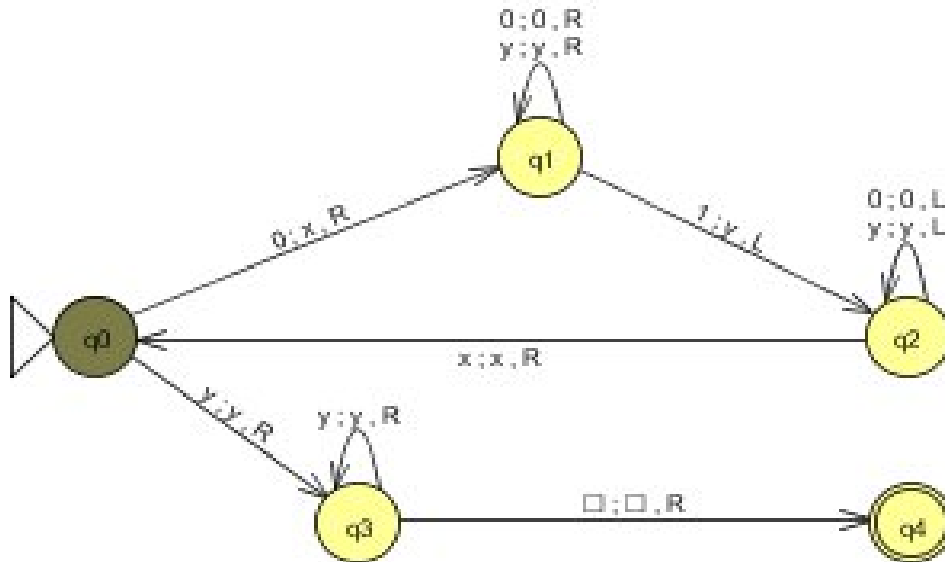
Step Reset Freeze Thaw Trace Remove





File Input Test Convert Help

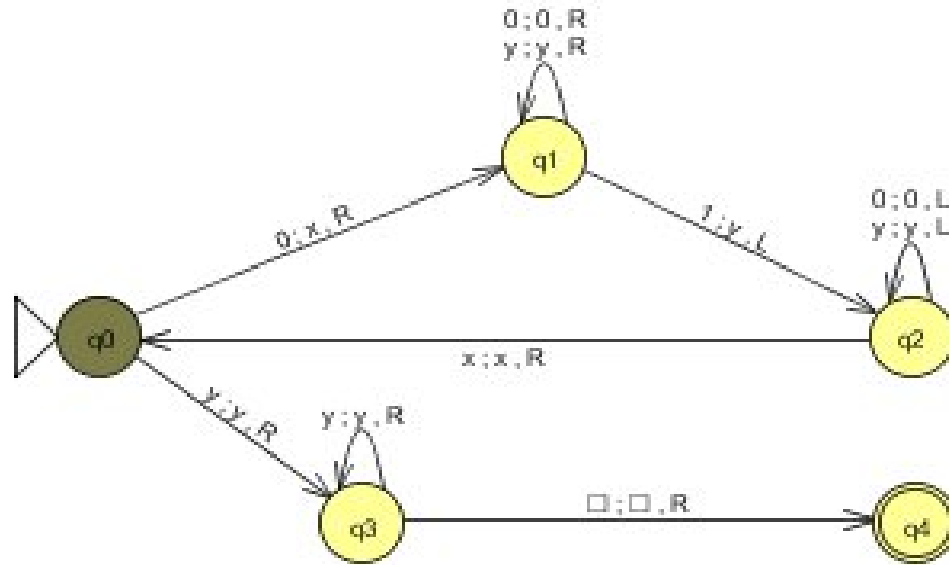
Editor Simulate: [10]



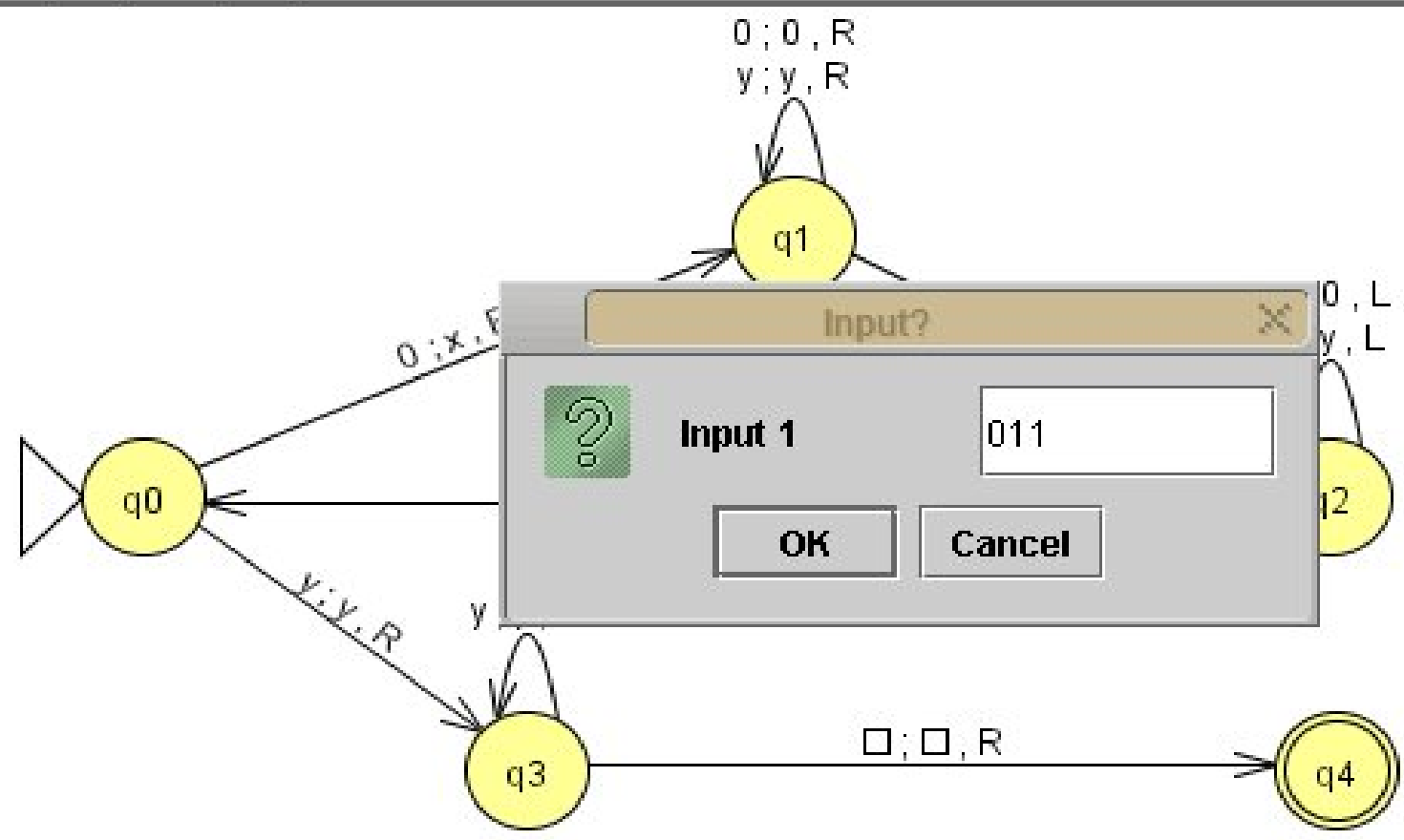
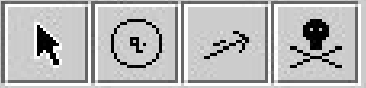
Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor Simulate: [10]



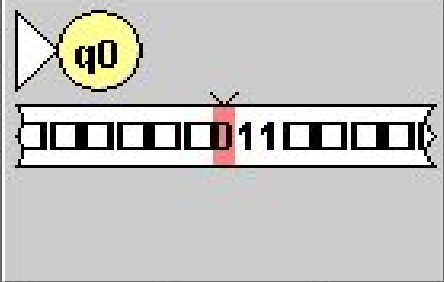
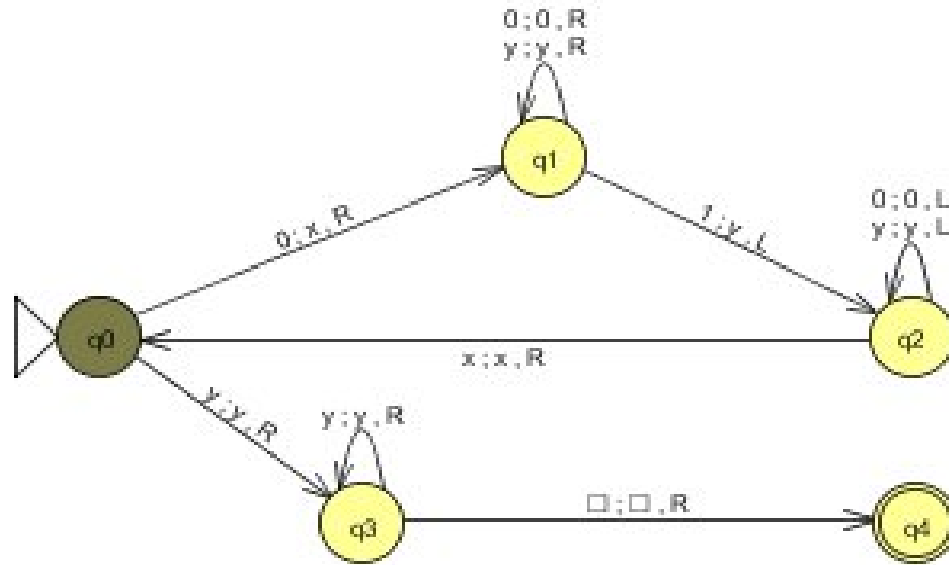
Step Reset Freeze Thaw Trace Remove



File Input Test Convert Help

Editor

Simulate: [011]

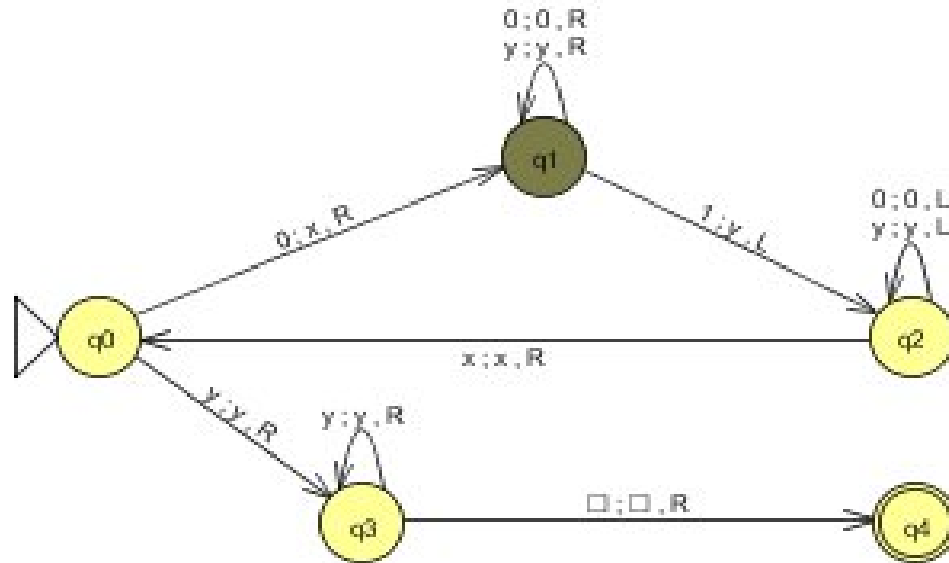


Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor

Simulate: [011]



q1



Step

Reset

Freeze

Thaw

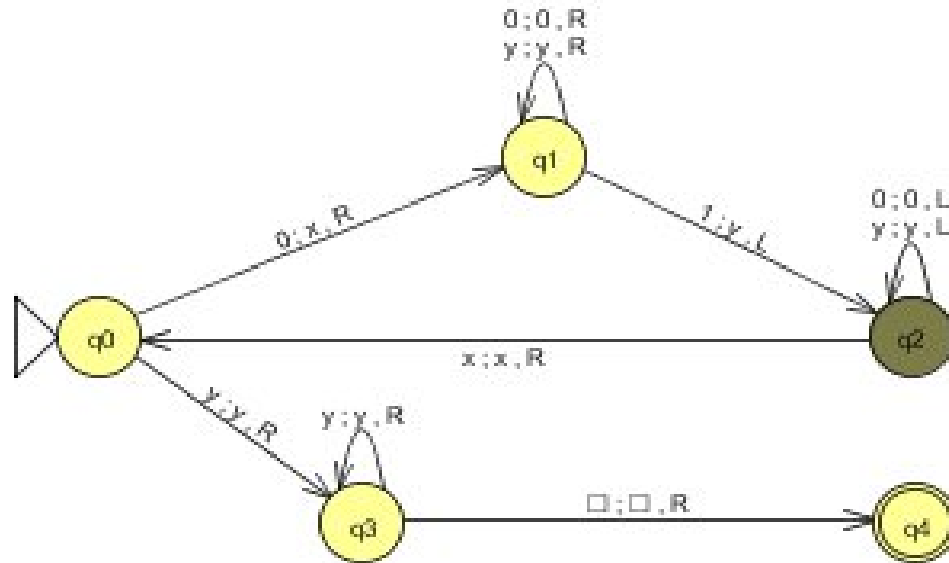
Trace

Remove

File Input Test Convert Help

Editor

Simulate: [011]



q2



Step

Reset

Freeze

Thaw

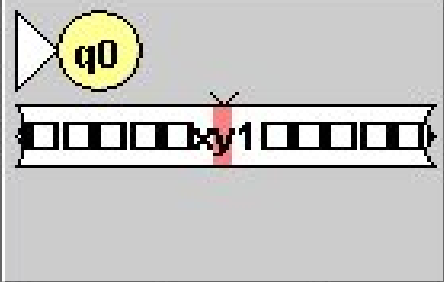
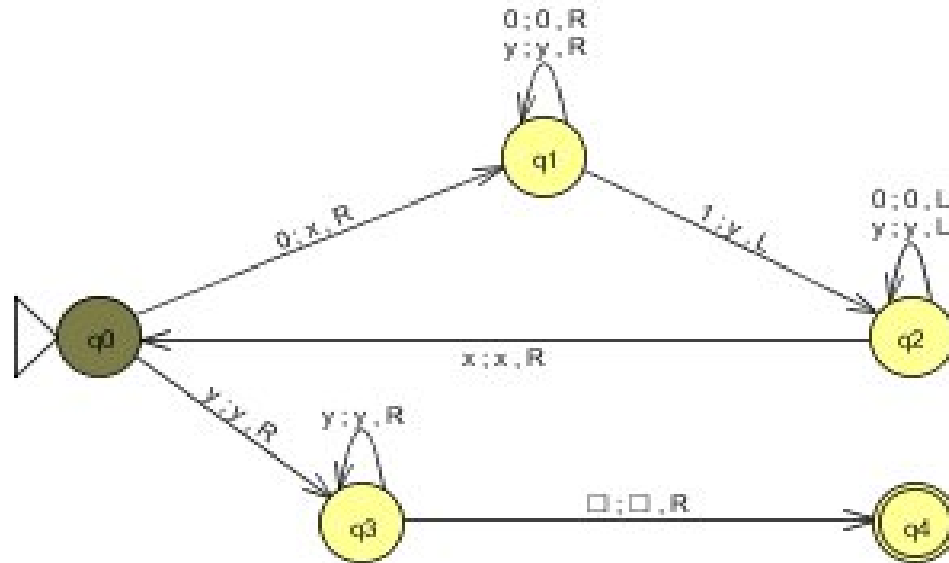
Trace

Remove

File Input Test Convert Help

Editor

Simulate: [011]

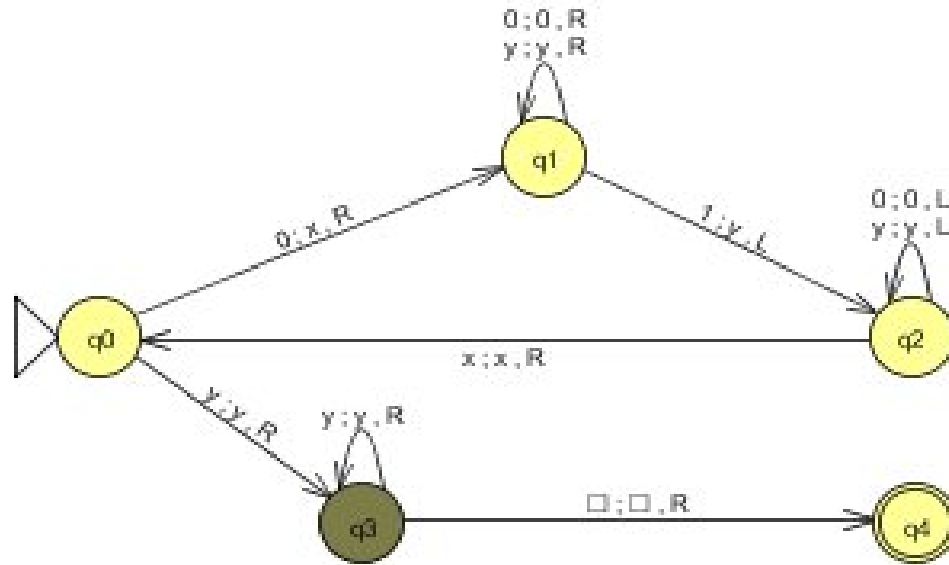


Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor

Simulate: [011]



q3



Step

Reset

Freeze

Thaw

Trace

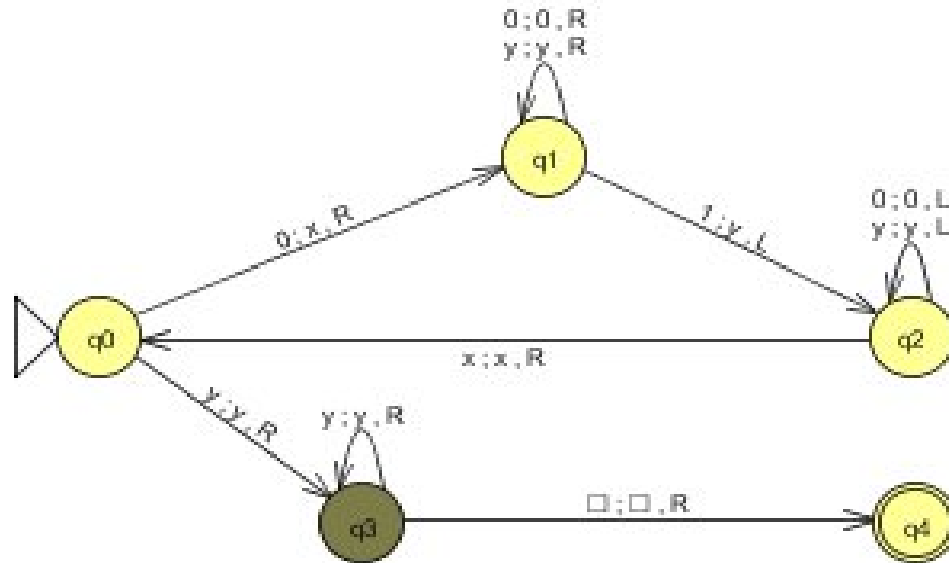
Remove



File Input Test Convert Help

Editor

Simulate: [011]



q3



Step

Reset

Freeze

Thaw

Trace

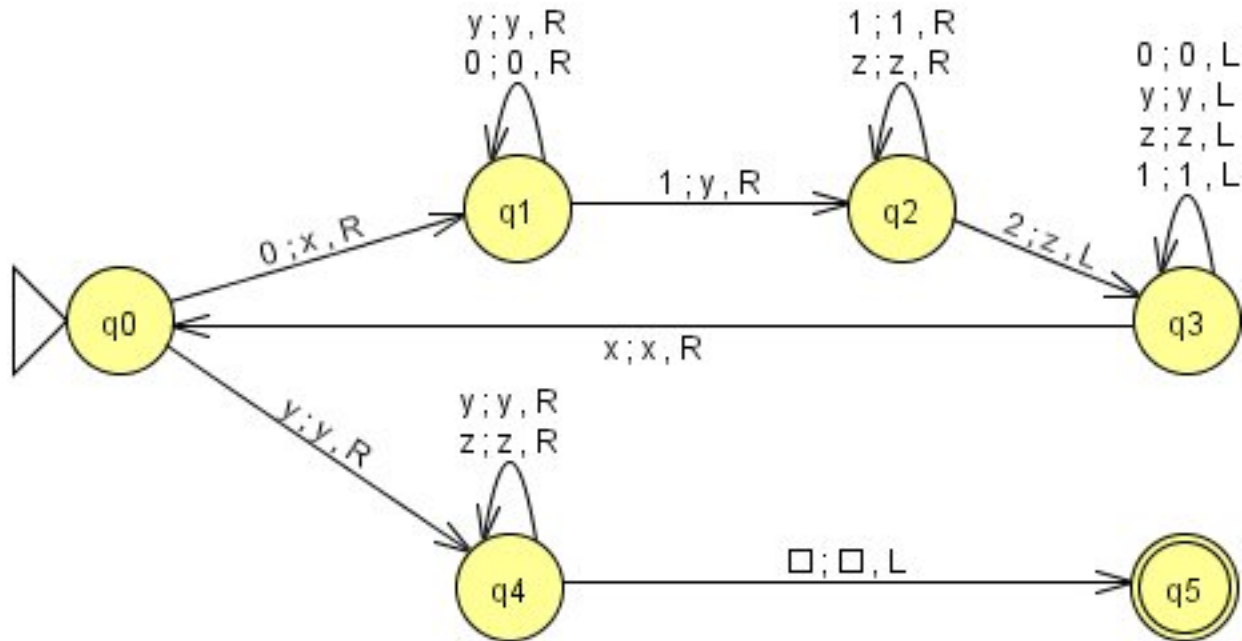
Remove

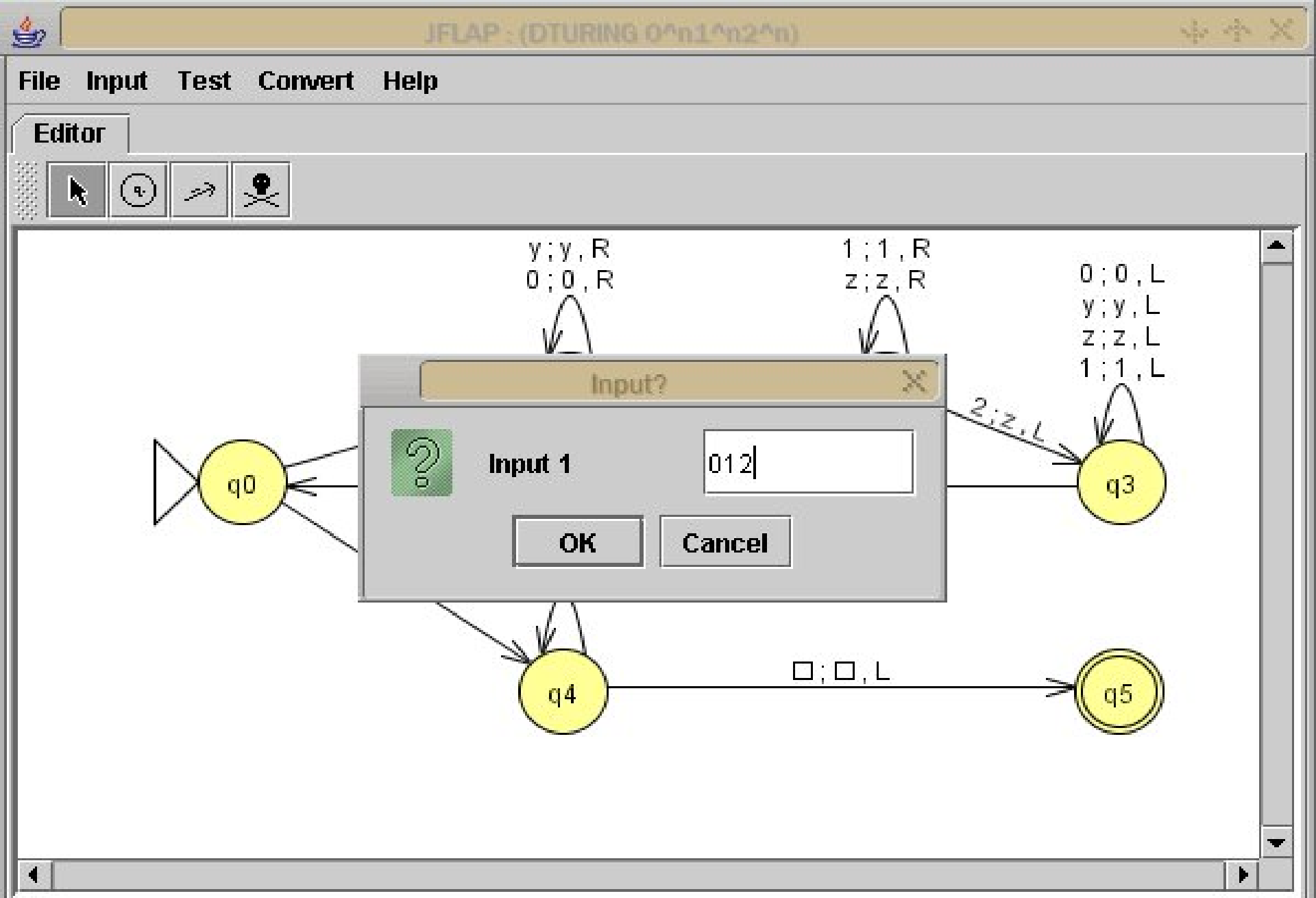
# TM for $\{0^n 1^n 2^n \mid n > 0\}$

M = “on input string w:

1. Look for 0's from the left end of the tape. Only blank cells are allowed to pass.
2. If 0 found, change it to x and move right, else reject
3. Scan to the right passing 0's and y's until reach 1
4. If 1 found, change it to y and move right, else reject.
5. Scan to the right passing 1's and z's until reach 2
6. If 2 found, change it to z and move left, else reject.
7. Scan to the left passing z's, 1's, y's, and 0's,
8. If x found move right
9. If 0 found, loop back to step 2.
10. If 0 not found, scan to the right passing only y's and z's and accept. Otherwise reject.

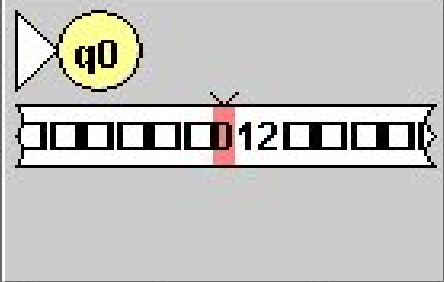
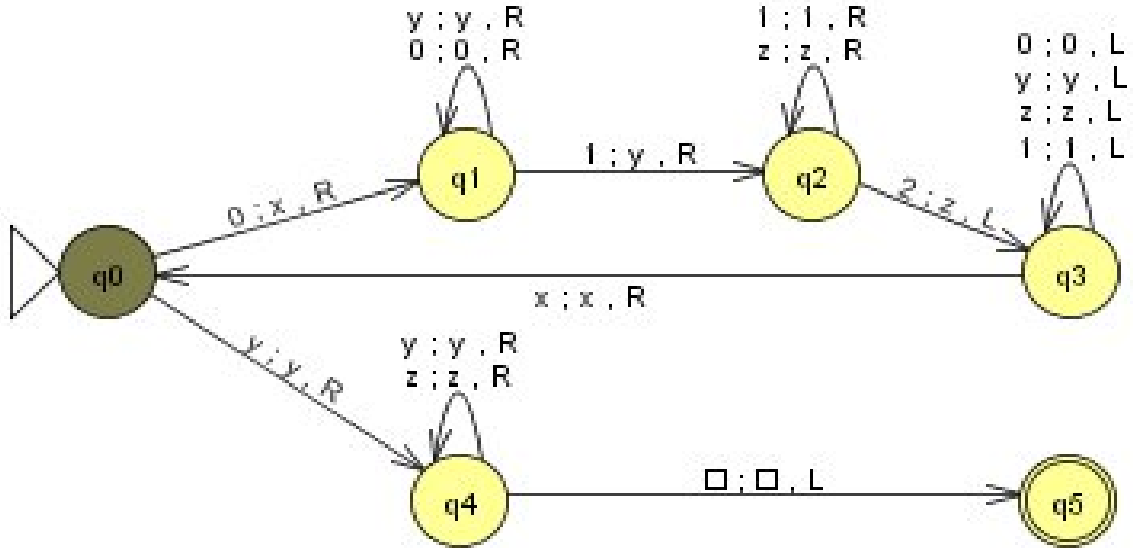
# Example of TM for $\{0^n 1^n 2^n \mid n > 0\}$ contd.





File Input Test Convert Help

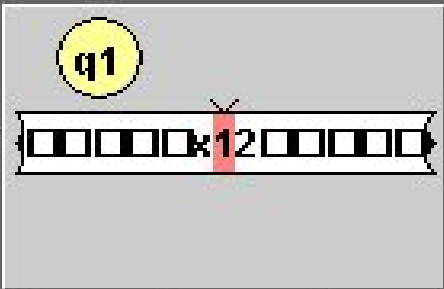
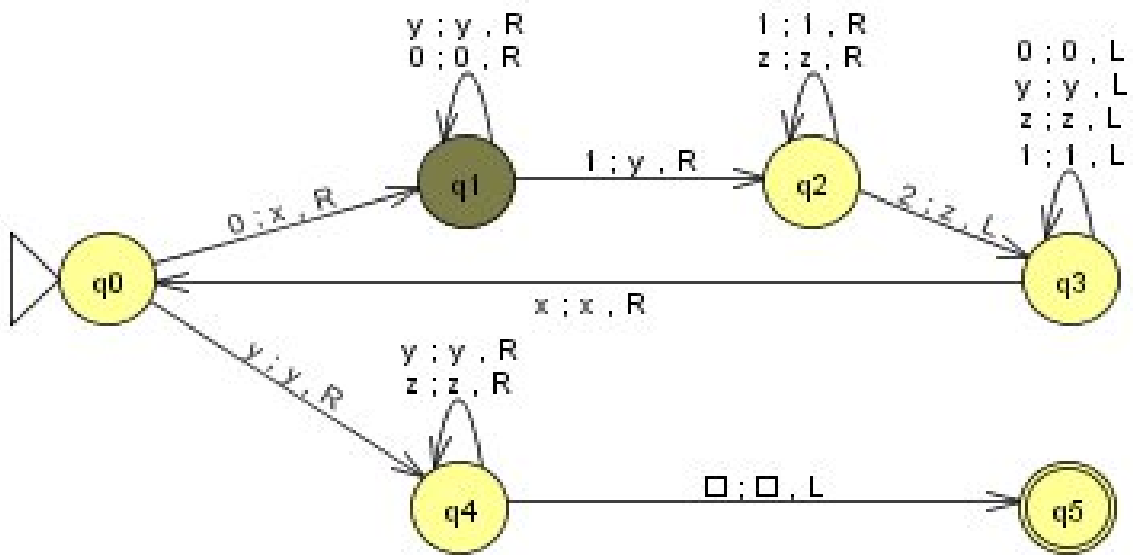
Editor Simulate: [012]



Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

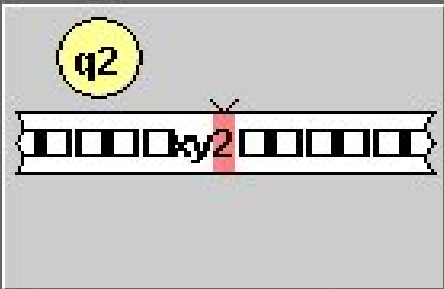
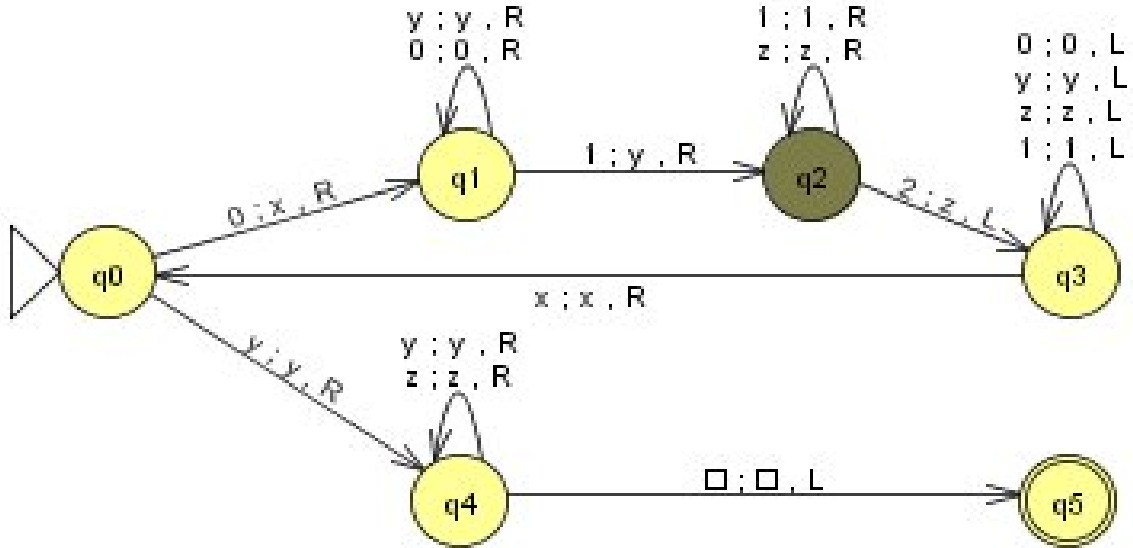
Editor Simulate: [012]



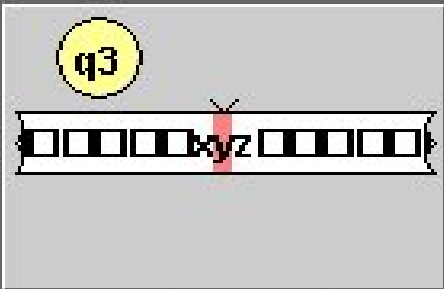
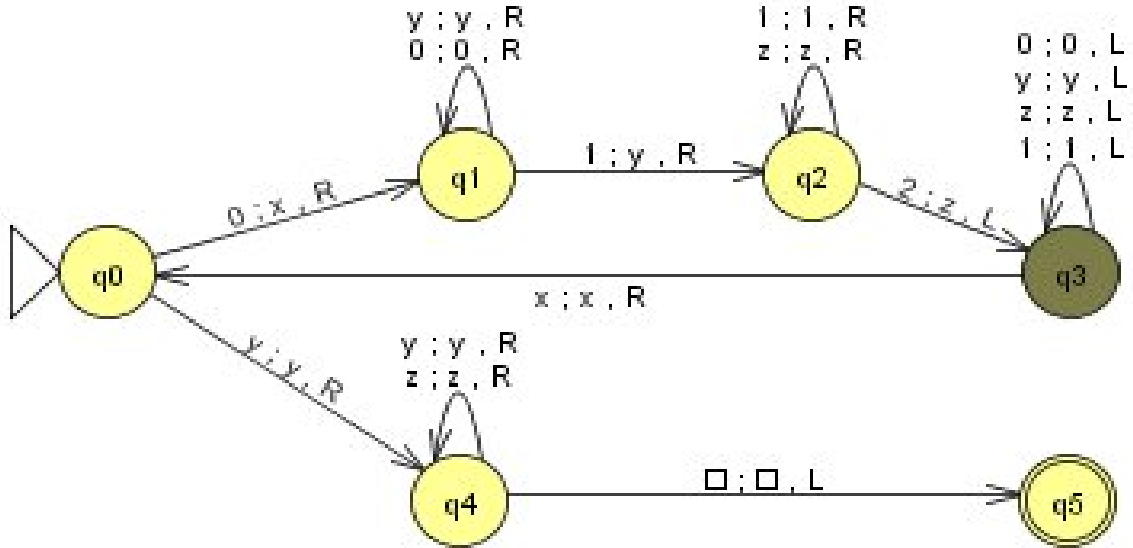
Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor Simulate: [012]



Step Reset Freeze Thaw Trace Remove

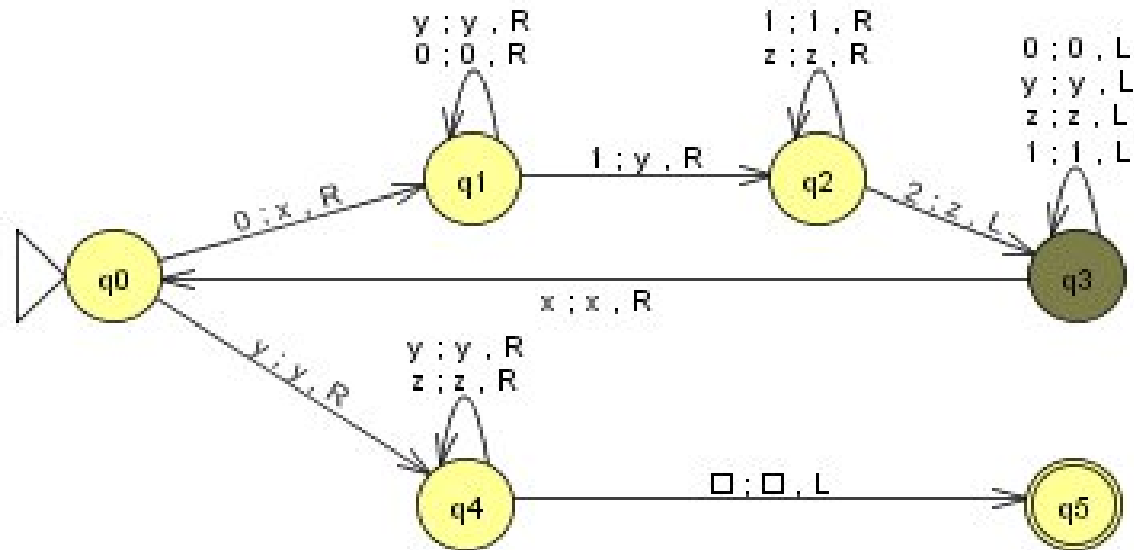




File Input Test Convert Help

Editor

Simulate: [012]



q3



Step

Reset

Freeze

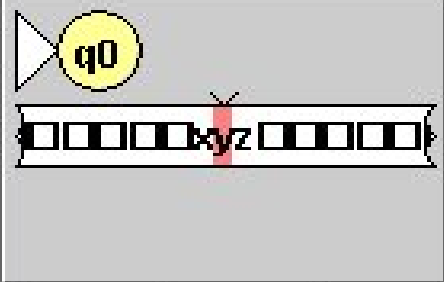
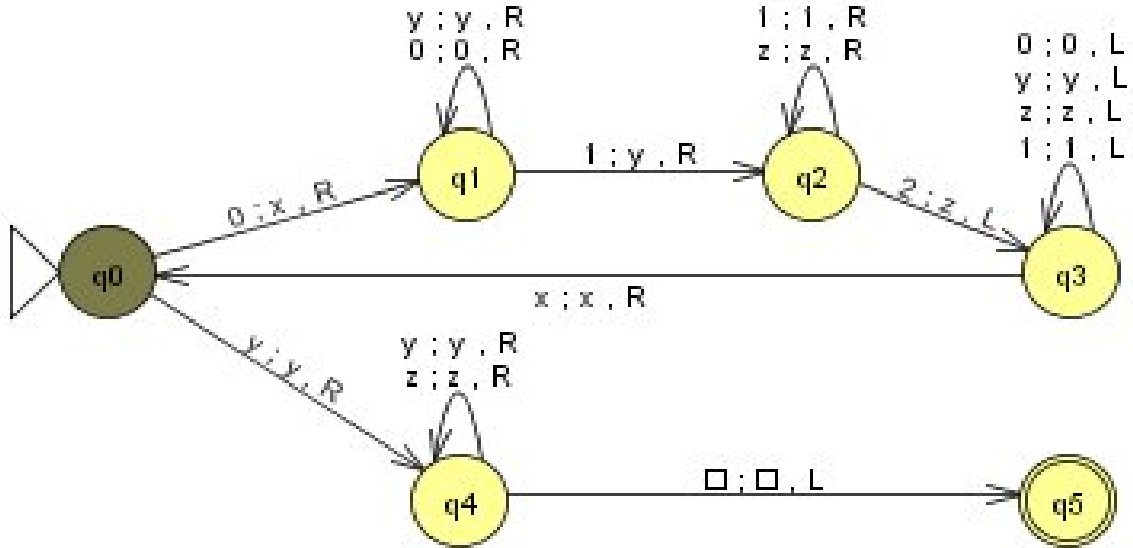
Thaw

Trace

Remove

File Input Test Convert Help

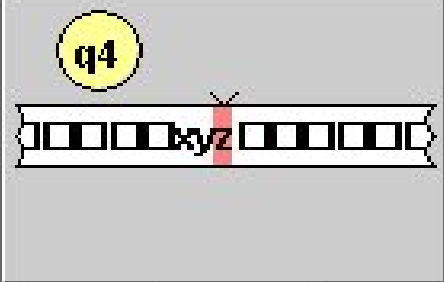
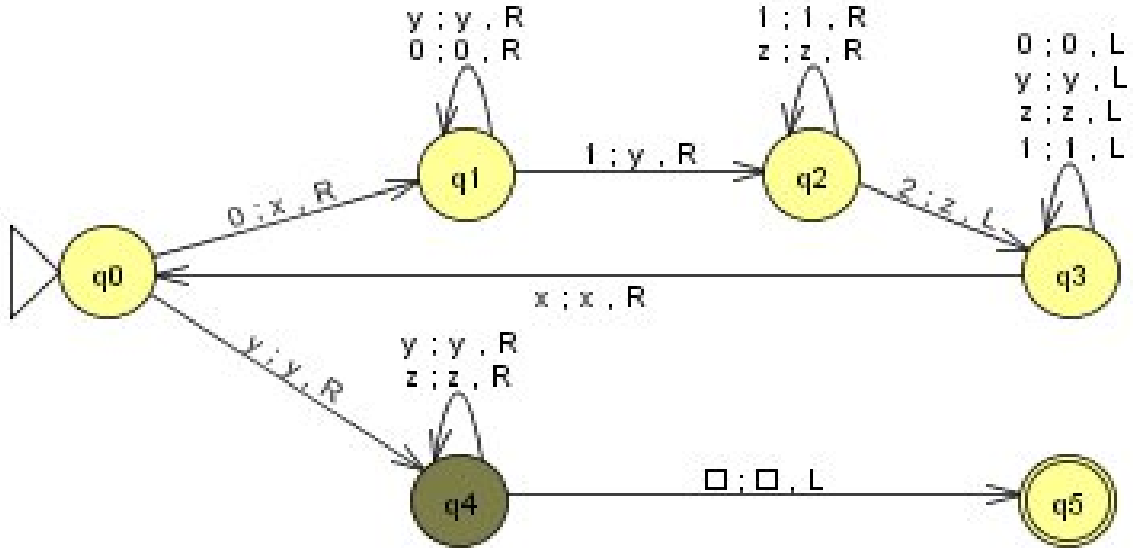
Editor Simulate: [012]



Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

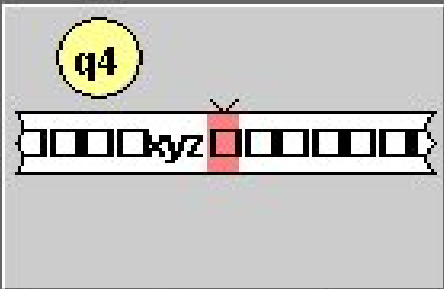
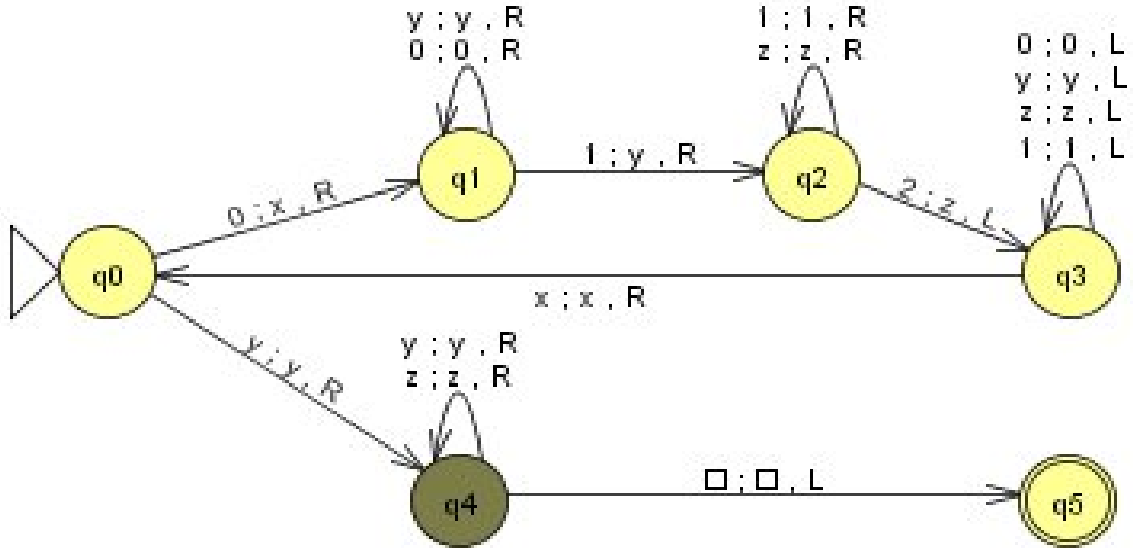
Editor Simulate: [012]



Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor Simulate: [012]

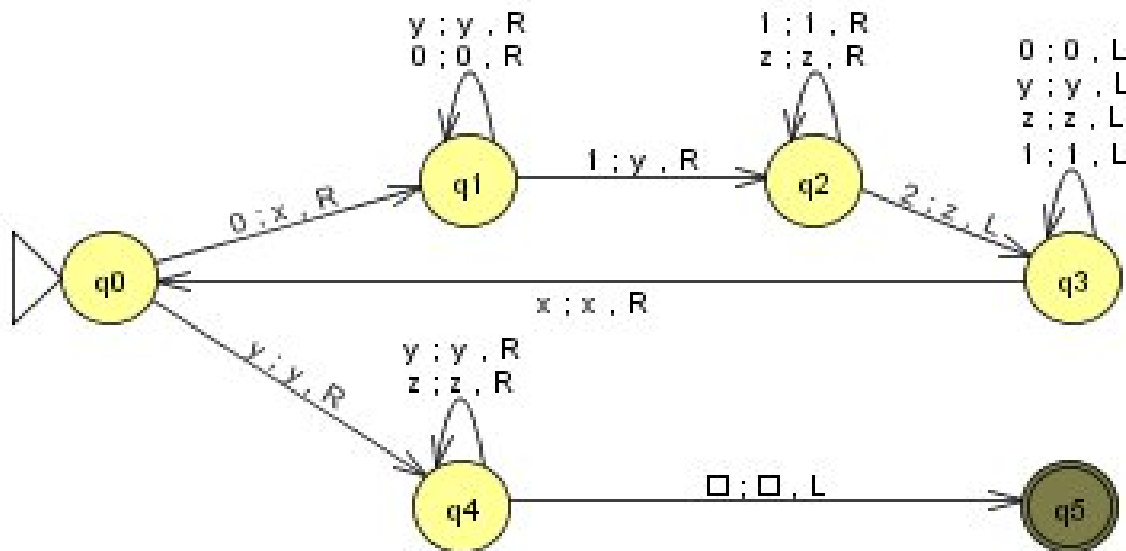


Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor

Simulate: [012]



Step

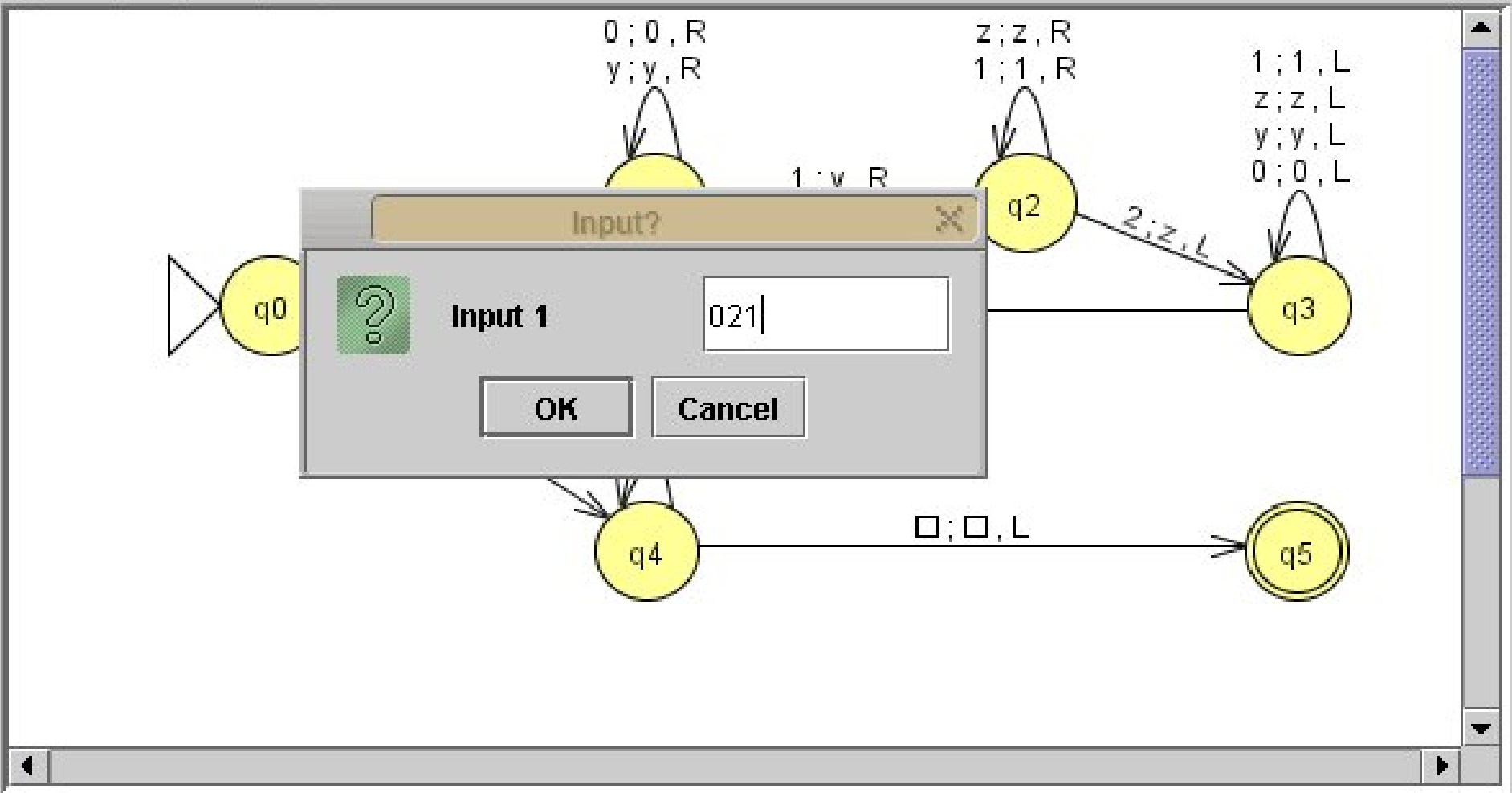
Reset

Freeze

Thaw

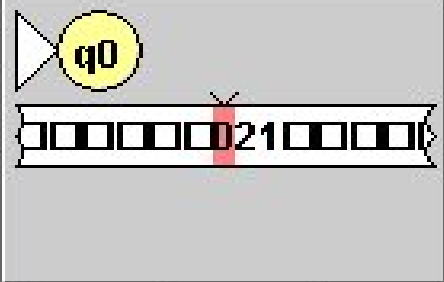
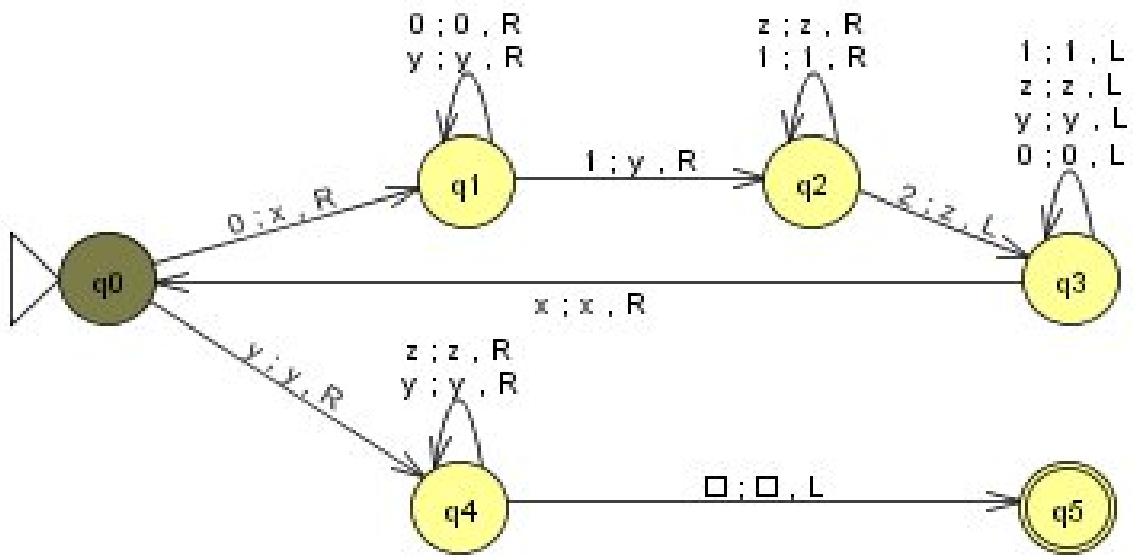
Trace

Remove



File Input Test Convert Help

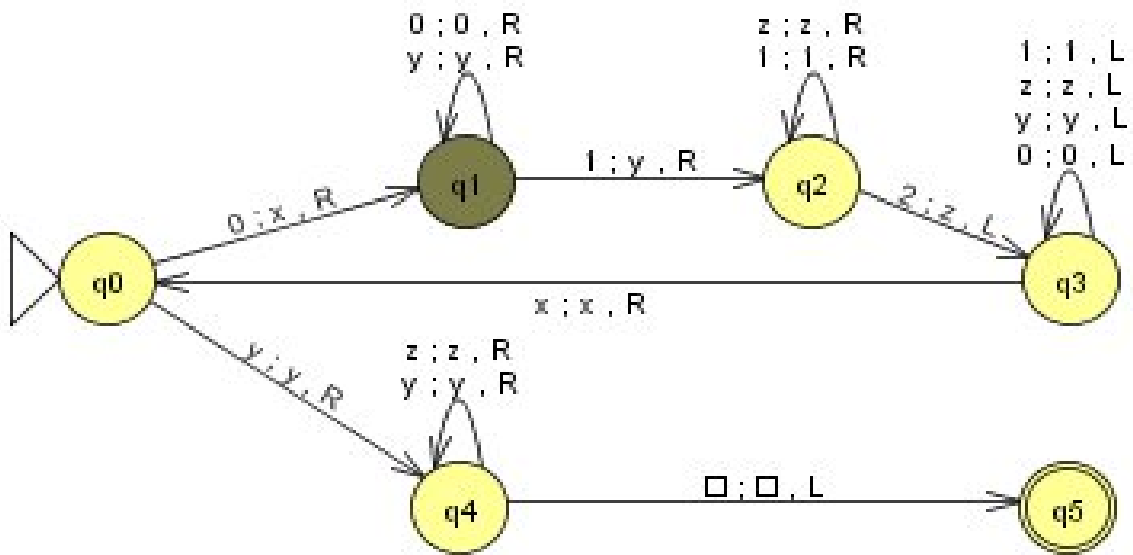
Editor Simulate: [021]



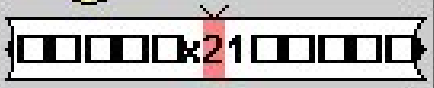
Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor Simulate: [021]



q1

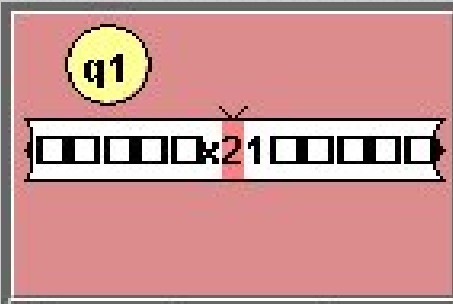
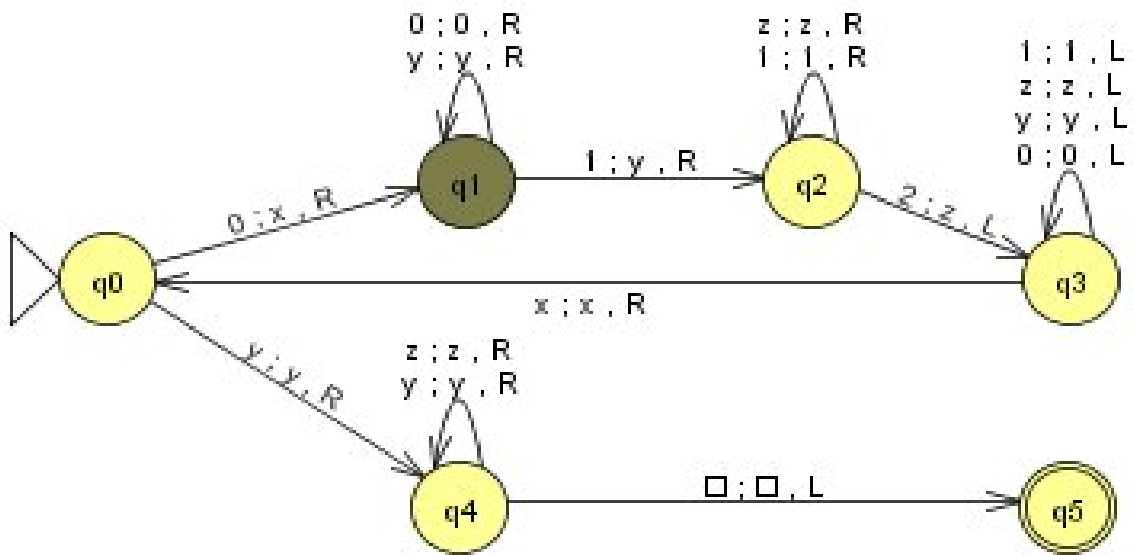


Step Reset Freeze Thaw Trace Remove

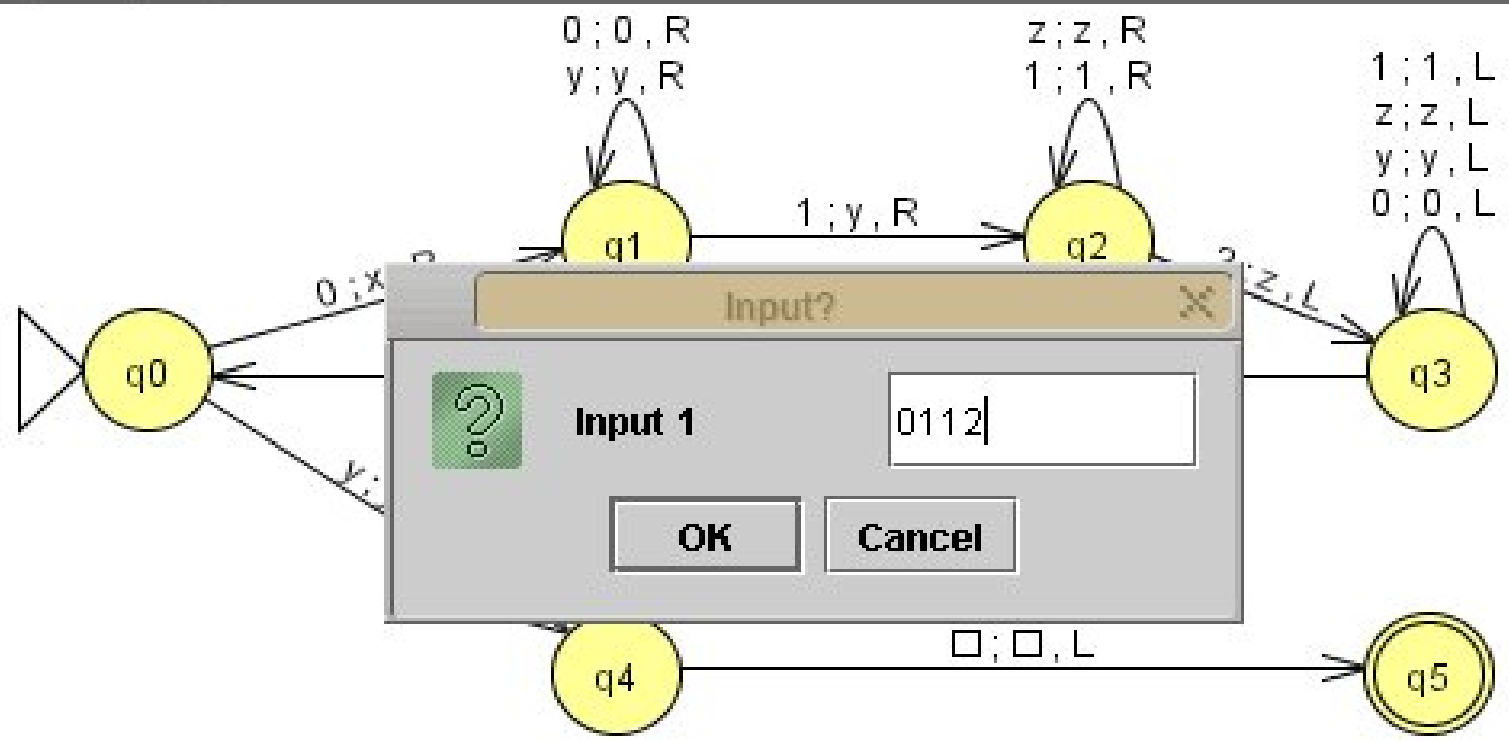


File Input Test Convert Help

Editor Simulate: [021]

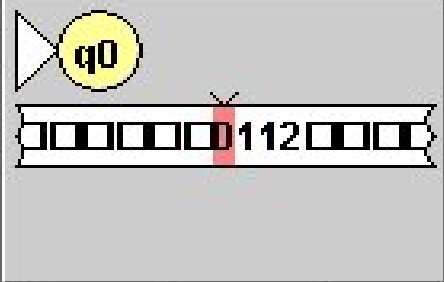
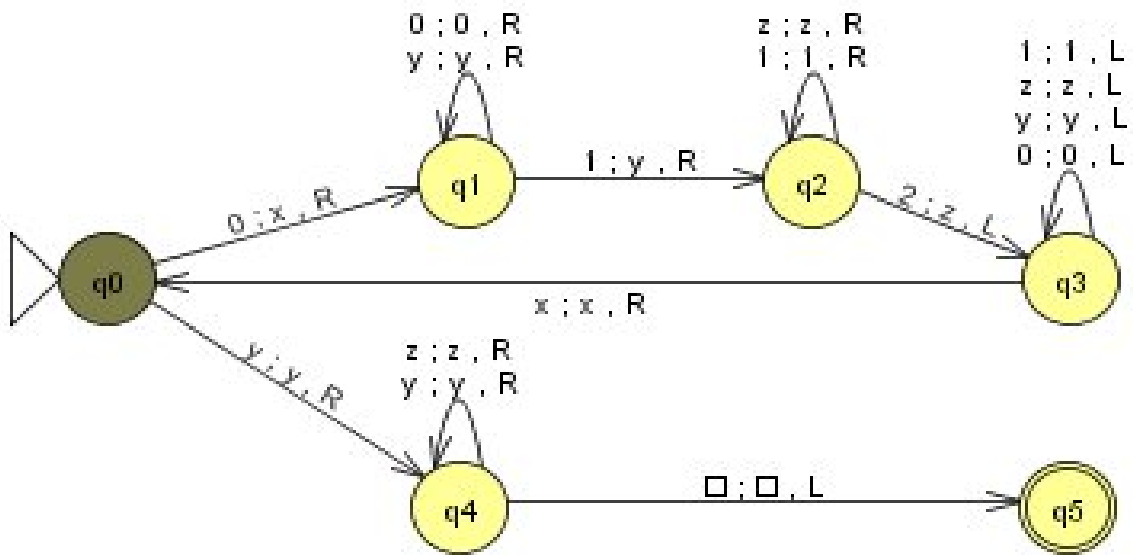


Step Reset Freeze Thaw Trace Remove



File Input Test Convert Help

Editor Simulate: [0112]

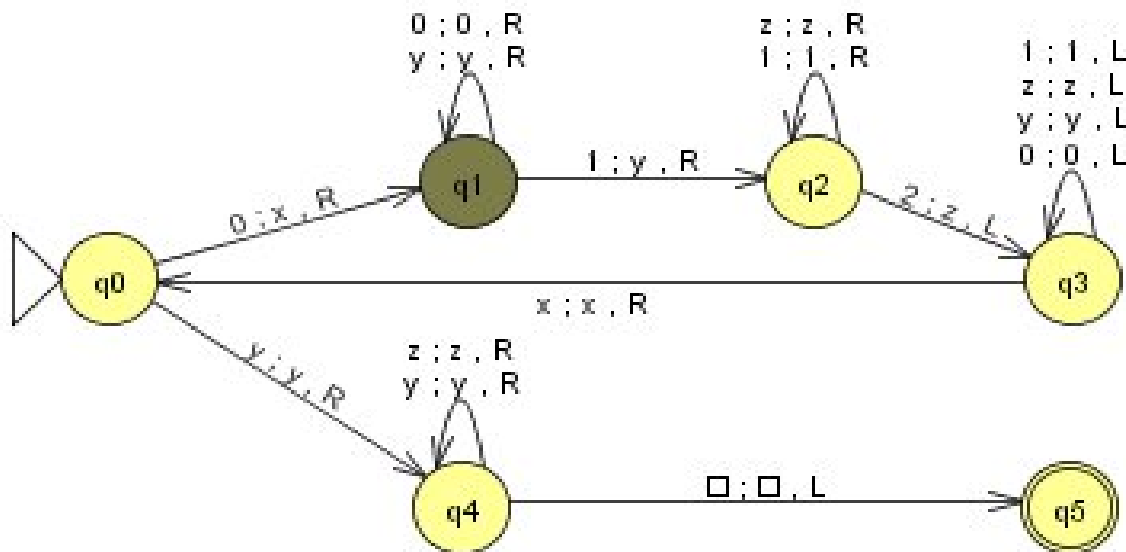


Step Reset Freeze Thaw Trace Remove

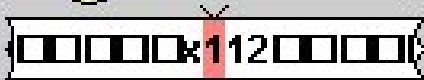
File Input Test Convert Help

Editor

Simulate: [0112]



q1



Step

Reset

Freeze

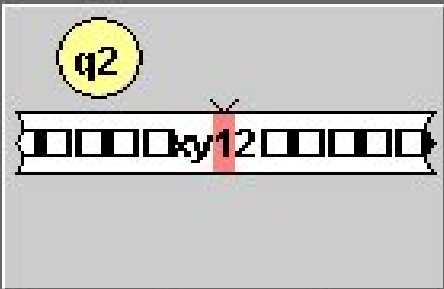
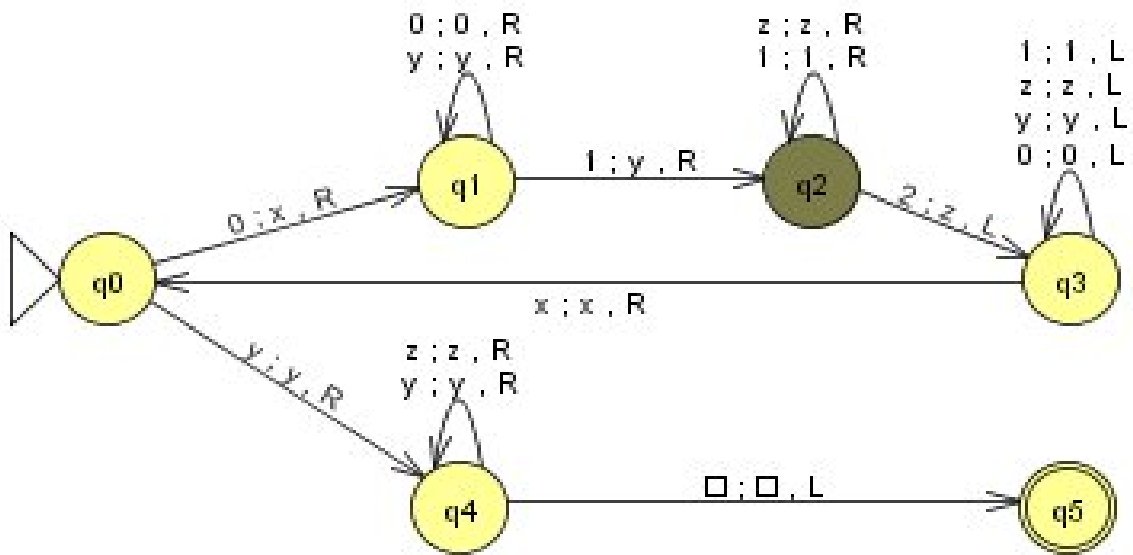
Thaw

Trace

Remove

File Input Test Convert Help

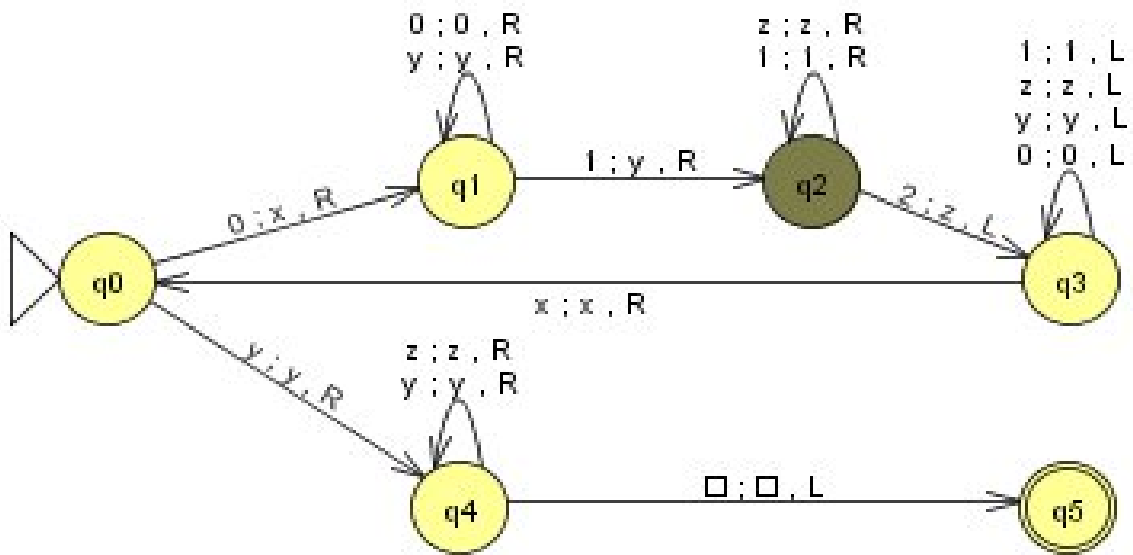
Editor Simulate: [0112]



Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

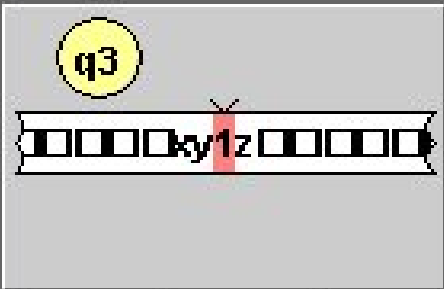
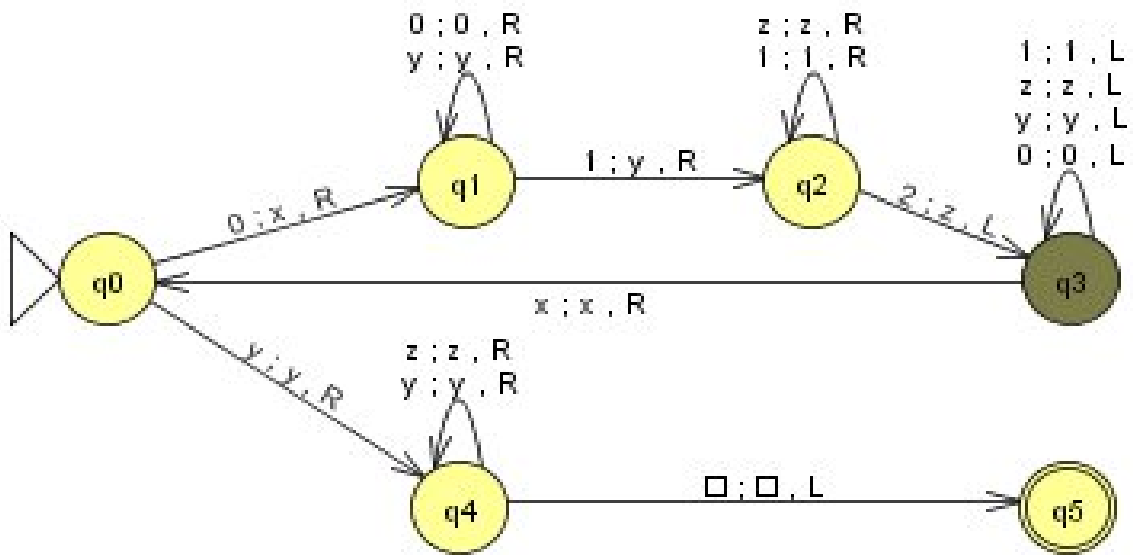
Editor Simulate: [0112]



Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

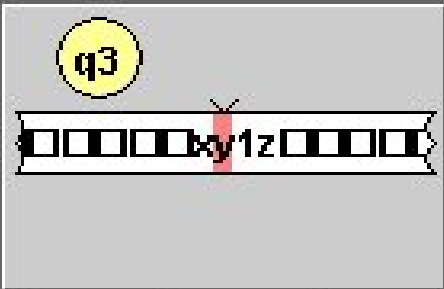
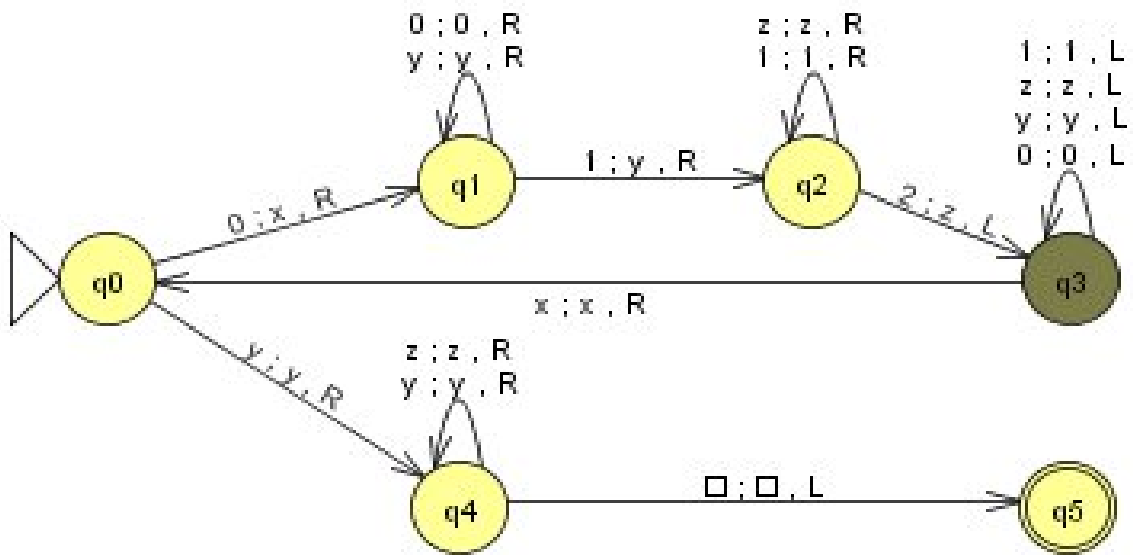
Editor Simulate: [0112]



Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

Editor Simulate: [0112]

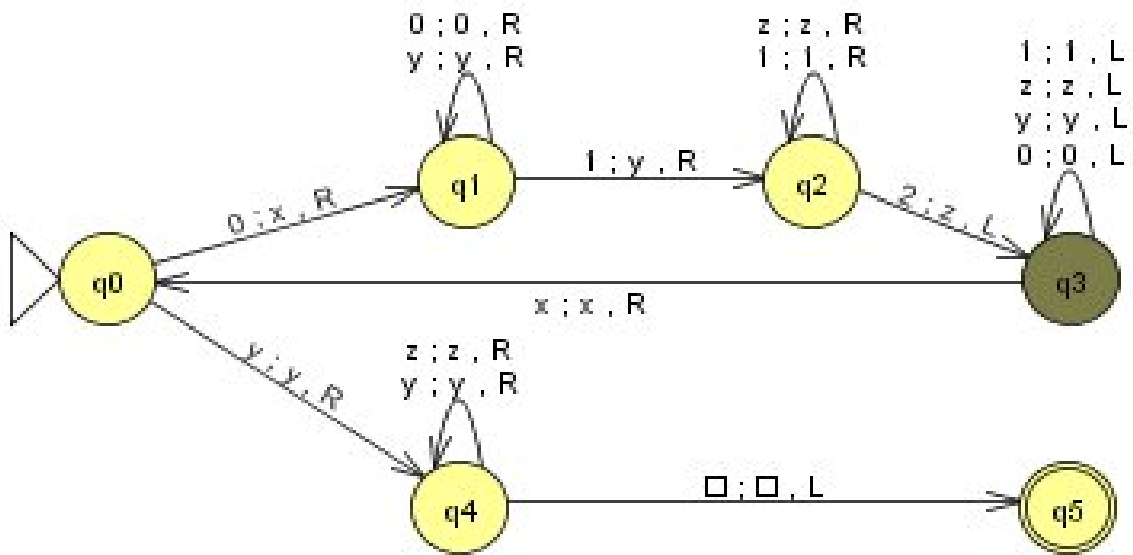


Step Reset Freeze Thaw Trace Remove



File Input Test Convert Help

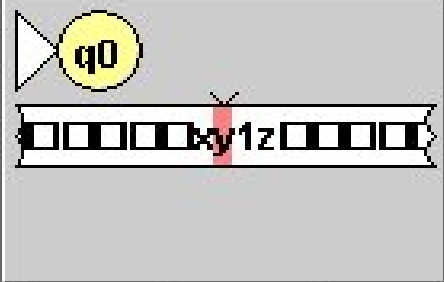
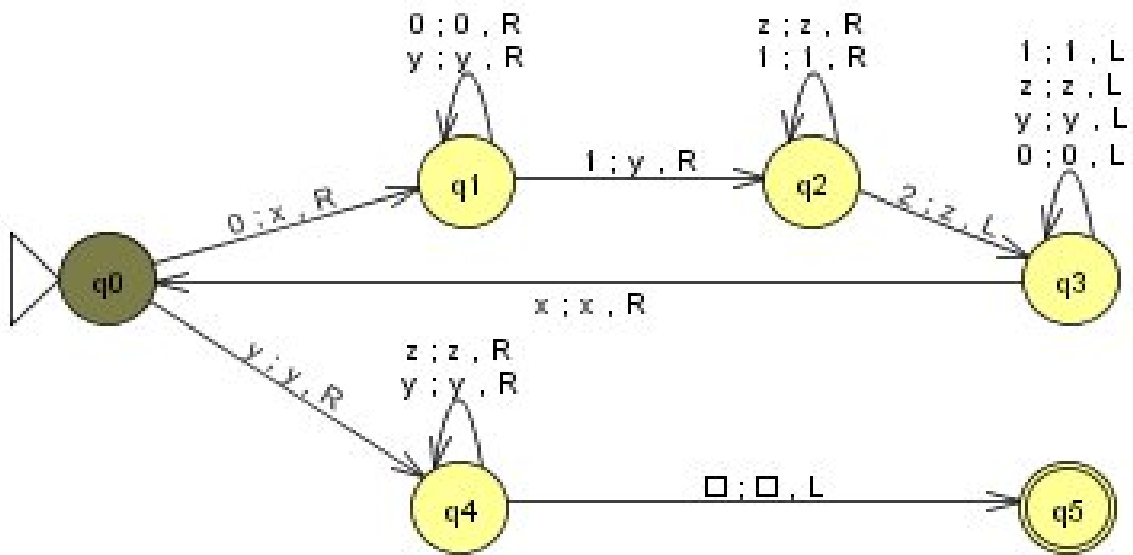
Editor Simulate: [0112]



Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

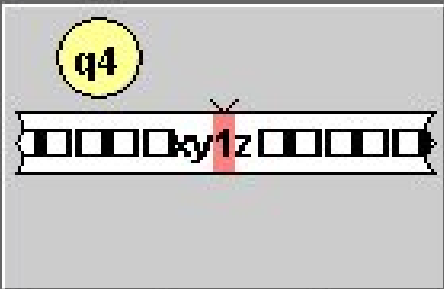
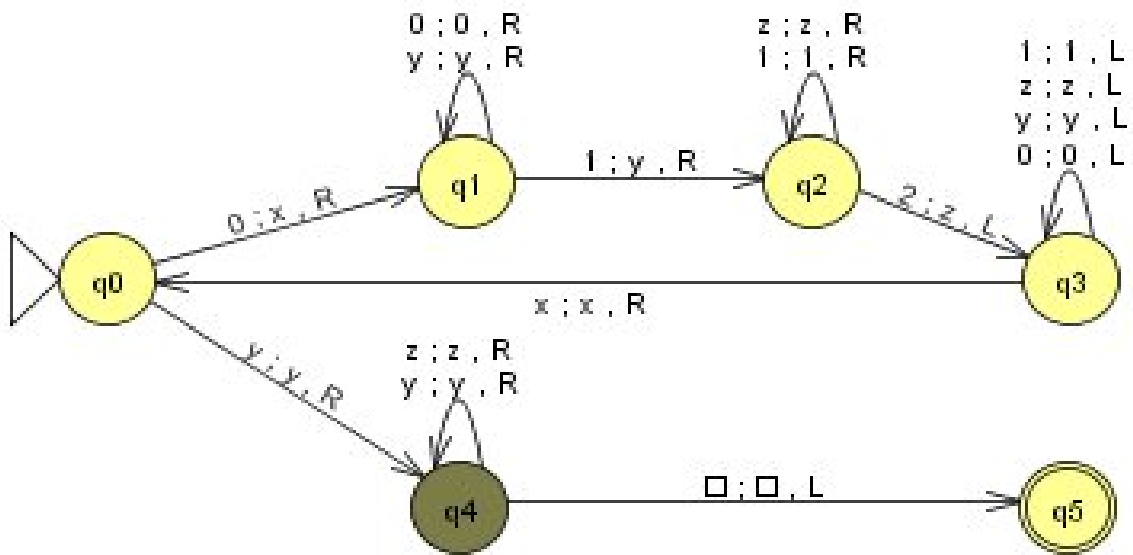
Editor Simulate: [0112]



Step Reset Freeze Thaw Trace Remove

File Input Test Convert Help

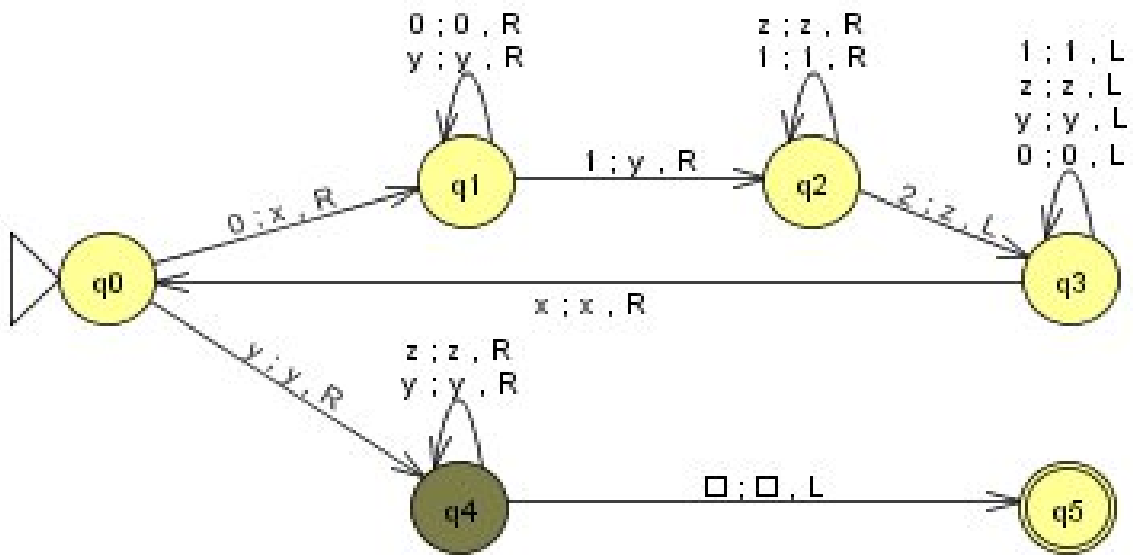
Editor Simulate: [0112]



Step Reset Freeze Thaw Trace Remove

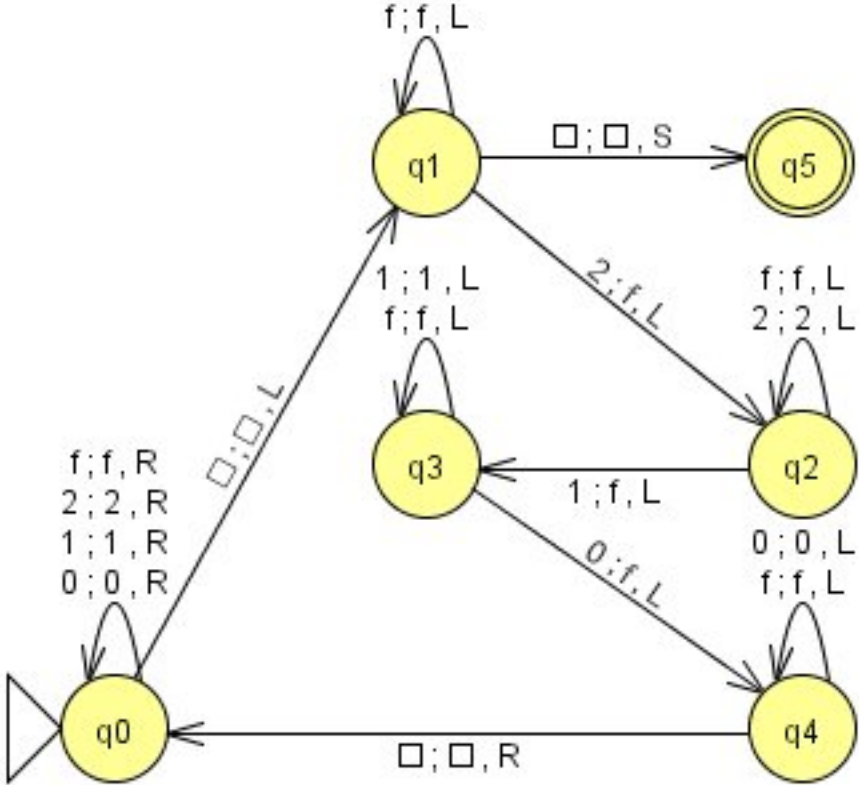
File Input Test Convert Help

Editor Simulate: [0112]



Step Reset Freeze Thaw Trace Remove

# Different way of making $\{0^n 1^n 2^n \mid n \geq 0\}$



# TM for $\{w\#w \mid w \in \{0, 1\}^*\}$

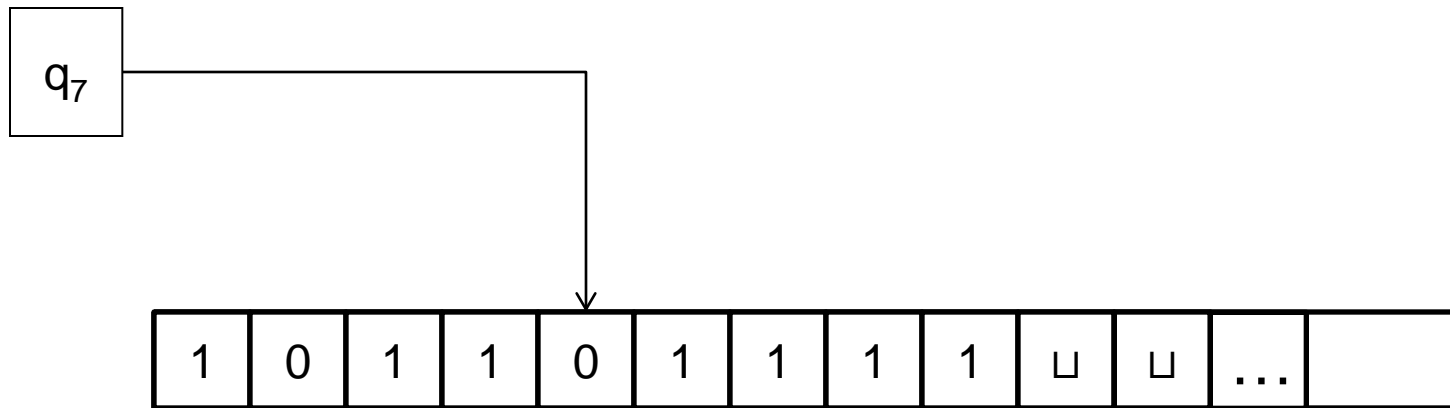
M = “ On input string  $w$ :

1. Check whether the string contains exactly one  $\#$ . If no, *reject*;
2. Check whether the string contains only 0's and 1's besides the  $\#$ . If there are other symbols, *reject*;
3. Zig-zag across the tape to corresponding positions on either side of the  $\#$  symbol to check whether these positions contain the same symbol. If they do not, *reject*;
4. Cross off symbols as they are checked to keep track of which symbols correspond;
5. When all symbols to the left of the  $\#$  have been crossed off, check for any remaining symbols to the right of the  $\#$ . If any symbols remain, *reject*; otherwise, *accept*.”

# Configurations

- A TM configuration is a setting of three items: current state, current tape contents, and the current head location, i.e.  $uqv$ , where  $u$  and  $v$  are substrings and  $q$  is a state.
- Configuration  $C_1$  yields configuration  $C_2$  if the Turing machine can legally go from  $C_1$  to  $C_2$  in a single step.
  - Suppose that  $a, b,$  and  $c$  in  $\Gamma$ ,  $u$  and  $v$  in  $\Gamma^*$  and states  $q_i$  and  $q_j$ .  $uaq_i b v$  and  $uq_j a c v$  are two configurations.
    - $uaq_i b v$  yields  $uq_j a c v$  if  $\delta(q_i, b) = (q_j, c, L)$ .
      - handles the case where the TM moves leftward.
    - For a rightward move,  $uaq_i b v$  yields  $uacq_j v$  if  $\delta(q_i, b) = (q_j, c, R)$ .

# Configurations (cont'd)



A Turing machine with configuration  $1011q_701111$



# Configurations: Special Cases

- Special Cases occur when the head is at the left-hand end or the right-hand end of the configurations.
  - For the left-hand end, the configurations  $q_i b v$  yields  $q_j c v$  if the transition is left-moving:  $\delta(q_i, b) = (q_j, c, L)$ .
  - For the right-hand end, the configuration  $u a q_i$  is equivalent to  $u a q_i \sqcup$  because we assume that blanks follow the part of the tape represented in the configuration. Thus we can handle this case as before, with the head no longer at the right-hand end.

# Configurations (cont'd)

- The start configuration of  $M$  on input  $w$  is the configuration  $q_0w$ , which indicates that the machine is in the start state  $q_0$  with its head at the leftmost position on the tape.
- In an accepting configuration, the state of the configuration is  $q_{\text{accept}}$ .
- In a rejecting configuration, the state of the configuration is  $q_{\text{reject}}$ .
- Accepting and rejecting configurations are halting configurations and do not yield further configurations.

# $L(M)$ : the language of TM $M$

- A TM  $M$  accepts input  $w$  if a sequence of configurations  $C_1, C_2, \dots, C_k$  exists, where
  1.  $C_1$  is the start configuration of  $M$  on input  $w$ ,
  2. Each  $C_i$  yields  $C_{i+1}$ , and
  3.  $C_k$  is an accepting configuration.
- The collection of strings that  $M$  accepts is the language of  $M$ , denoted  $L(M)$ .

# Turing decidable/recognizable

- Turing-recognizable
  - A language is Turing-recognizable if some Turing machine recognizes it.
  - i.e., TM  $M$  recognizes language  $L$  if  $L = \{w \mid M \text{ accepts } w\}$ .
  - *Note:* 3 outcomes possible, either TM accepts, rejects, or loops on a string.

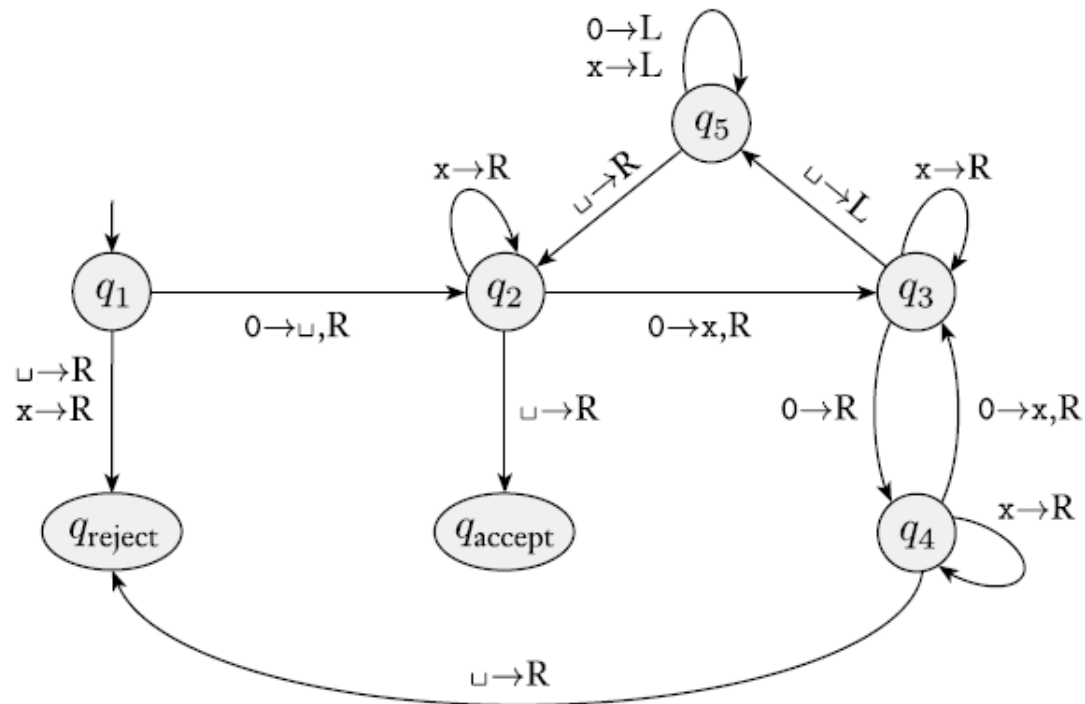
# Turing decidable/recognizable (cont'd)

- Turing-decidable:
  - A language is Turing-decidable if some Turing machine decides it.
  - TM  $M$  decides  $L$  if
    - (i)  $w \in L$ ,  $M$  accepts it.
    - (ii)  $w \notin L$ ,  $M$  rejects it.

Note: Every decidable language is Turing-recognizable but certain Turing-recognizable language are not decidable.

# TM for $\{0^{2^n} \mid n \geq 0\}$

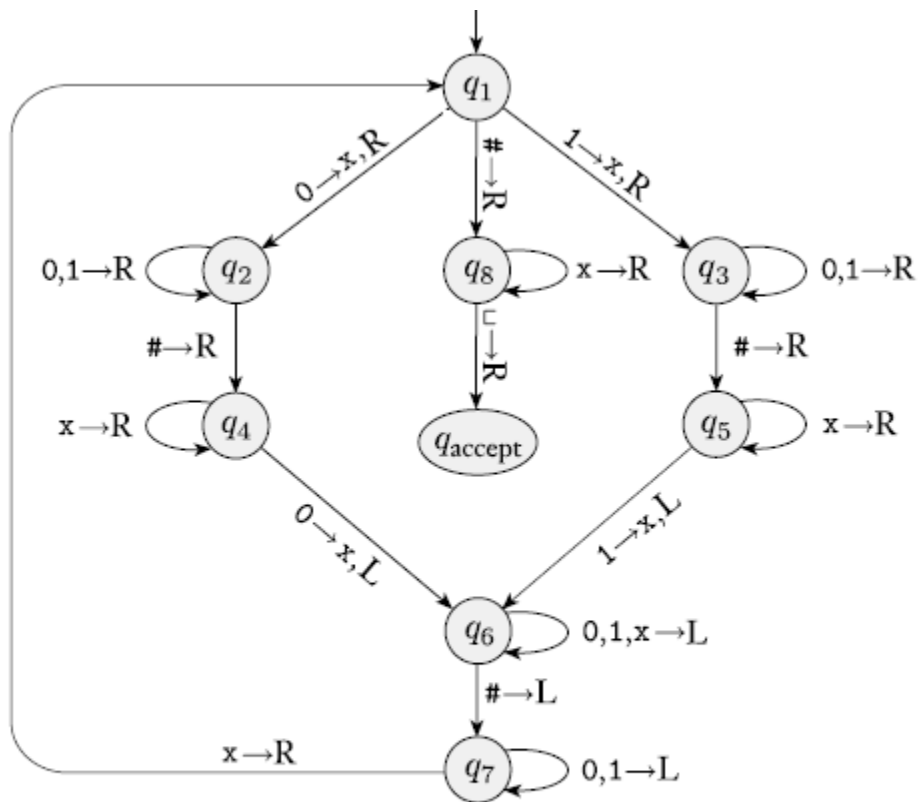
- TM M= “on input string  $w$ :
  1. Sweep left to right across the tape, crossing off every other 0,
  2. If in step 1, the tape contained a single 0, accept,
  3. If in step 1, the tape contained more than a single 0 and the number of 0s was odd, reject,
  4. Return the head to the left-hand end of the tape,
  5. Go to step 1.”



**FIGURE 3.8**  
State diagram for Turing machine  $M_2$

$q_1 0000$	$\sqcup q_5 x0x\sqcup$	$\sqcup xq_5 xx\sqcup$
$\sqcup q_2 000$	$q_5 \sqcup x0x\sqcup$	$\sqcup q_5 xxx\sqcup$
$\sqcup xq_3 00$	$\sqcup q_2 x0x\sqcup$	$q_5 \sqcup xxx\sqcup$
$\sqcup x0q_4 0$	$\sqcup xq_2 0x\sqcup$	$\sqcup q_2 xxx\sqcup$
$\sqcup x0xq_3 \sqcup$	$\sqcup xxxq_3 x\sqcup$	$\sqcup xq_2 xx\sqcup$
$\sqcup x0q_5 x\sqcup$	$\sqcup xxxq_3 \sqcup$	$\sqcup xxq_2 x\sqcup$
$\sqcup xq_5 0x\sqcup$	$\sqcup xxxq_5 x\sqcup$	$\sqcup xxxq_2 \sqcup$
		$\sqcup xxx\sqcup q_{\text{accept}}$

The figures are taken from the book *Introduction to Theory of Computation*, Michael Sipser, page 172.



**FIGURE 3.10**  
State diagram for Turing machine  $M_1$

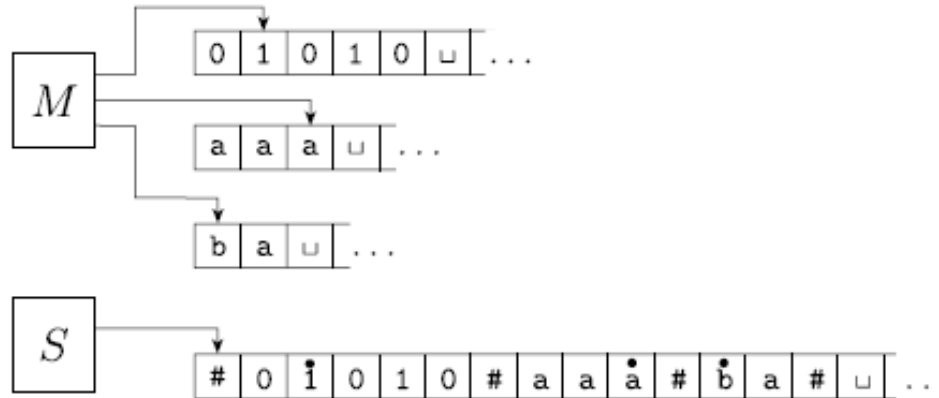
The figures are taken from the book *Introduction to Theory of Computation*, Michael Sipser, page 173.



# Variants of TM models

- Variants of TM model (not more powerful than basic TM model)
  - Multitape Turing machine
    - Like ordinary TM but with several tapes.
    - Every multitape Turing machine has an equivalent single-tape Turing machine

# Variants of TM models (cont'd)

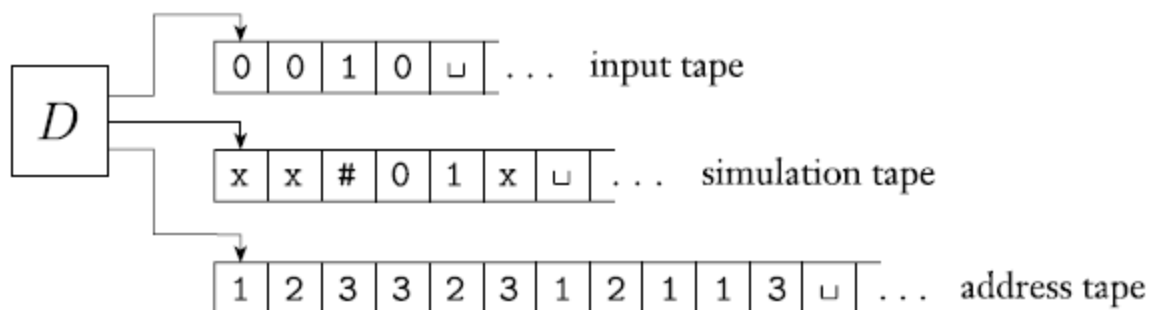


**FIGURE 3.14**  
Representing three tapes with one

This figure is taken from the book *Introduction to Theory of Computation*, Michael Sipser, page 177.

# Variants of TM models (cont'd)

- Nondeterministic Turing machine
  - Every nondeterministic Turing machine has a deterministic Turing machine



**FIGURE 3.17**  
Deterministic TM  $D$  simulating nondeterministic TM  $N$

This figure is taken from the book *Introduction to Theory of Computation*, Michael Sipser, page 179.

# Acknowledgements and Reference

- Thanks Dr. Rakesh Verma in University of Houston for his permission to using his slides for his course COSC 3340 (those with footnote indicted).
- <http://www2.cs.uh.edu/~rmverma/3340/3340.html>